

Relations Between Probability Measures

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1 Definitions

In this TED talk, we will talk about the relation between Kullback-Leibler divergence and the total variation distance between probability distributions. We will focus on **discrete** probability distributions only. Let's start with some definitions and assumed theorems first:

Definition 1.1. S is a sample space, $\mathcal{A} = 2^S$. A probability distribution $P : \mathcal{A} \rightarrow [0, 1]$ is a probability distribution. P satisfies Kolmogorov axioms of probability.

For $s \in S$, we also use the shorthand notation $P(s) := P(\{s\})$.

Definition 1.2. The Kullback-Leibler divergence between two probability distributions $P(x)$ and $Q(x)$ from discrete probability spaces defined over the same S is

$$D_{\text{KL}}(P||Q) = \sum_{x \in S} P(x) \log \frac{P(x)}{Q(x)}. \quad (1)$$

Definition 1.3. The Manhattan distance (L_1 metric) between two probability distributions $P(x)$ and $Q(x)$ from discrete probability spaces defined over the same S is

$$\|P - Q\|_1 = \sum_{x \in S} |P(x) - Q(x)|. \quad (2)$$

Definition 1.4. The total variation distance between two probability distributions $P(x)$ and $Q(x)$ from discrete probability spaces defined over the same S is

$$\delta(P, Q) = \max_{A \in 2^S} |P(A) - Q(A)|. \quad (3)$$

Theorem 1.5. Jensen's inequality.

For a convex function f , and reals $p_1, \dots, p_n \geq 0$ such that $\sum_{i=1}^n p_i = 1$ it holds that:

$$f\left(\sum_{i=1}^n p_i x_i\right) \leq \sum_{i=1}^n p_i f(x_i) \quad (4)$$

2 Lower Bound

Theorem 2.1. Pinsker's inequality. For two probability distributions $P(x)$ and $Q(x)$ from discrete probability spaces defined over the same S , it holds that

$$\|P - Q\|_1 \leq \sqrt{2 D_{\text{KL}}(P||Q)}.$$

The equivalent inequality is that

$$D_{\text{KL}}(P||Q) \geq \frac{1}{2} \|P - Q\|_1^2.$$

Proof. Bernoulli distributions case.

Let's denote by P and Q Bernoulli distribution over $S = \{0, 1\}$. Also, denote:

$$\begin{aligned} p &= P(0), 1 - p = P(1) \\ q &= Q(0), 1 - q = Q(1) \end{aligned}$$

We can see that

$$\begin{aligned}\|P - Q\|_1 &= |p - q| + |1 - p - 1 + q| = 2 \cdot |p - q| \\ \|P - Q\|_1^2 &= 4(p - q)^2\end{aligned}$$

Let's define $f(p, q) = D_{\text{KL}}(P||Q) - \frac{1}{2}\|P - Q\|_1^2$. We will analyse the behaviour of the function using basic calculus.

$$\begin{aligned}f(p, q) &= p \log \frac{p}{q} + (1 - p) \log \frac{1 - p}{1 - q} - 2(p - q)^2 \\ \frac{\partial f}{\partial q} &= p \cdot \frac{q}{p} \cdot \frac{-p}{q^2} + (1 - p) \cdot \frac{1 - q}{1 - p} \cdot (1 - p) \frac{-1}{(1 - q)^2} \cdot (-1) - 4(p - q) \cdot (-1) \\ &= -\frac{p}{q} + \frac{1 - p}{1 - q} + 4(p - q) \\ &= \frac{-p + pq + q - pq}{q(1 - q)} - 4(q - p) \\ &= (q - p) \left[\frac{1}{q(1 - q)} - 4 \right]\end{aligned}$$

We see that with $q \neq \frac{1}{2}$, the sign of the partial derivative depends only on the sign of $q - p$. Therefore, $\frac{\partial f}{\partial q}$ is negative for $q < p$, positive for $q > p$ and 0 for $q = p$. That means that for $q \neq \frac{1}{2}$, $q = p$ is the minimum of $f(p, q)$.

$$\begin{aligned}f(p, p) &= p \log \frac{p}{p} + (1 - p) \log \frac{1 - p}{1 - p} - 2(p - p)^2 \\ &= 0.\end{aligned}$$

That means that for $q \neq \frac{1}{2}$, $f(p, q)$ is non-negative and the Pinsker's inequality holds.

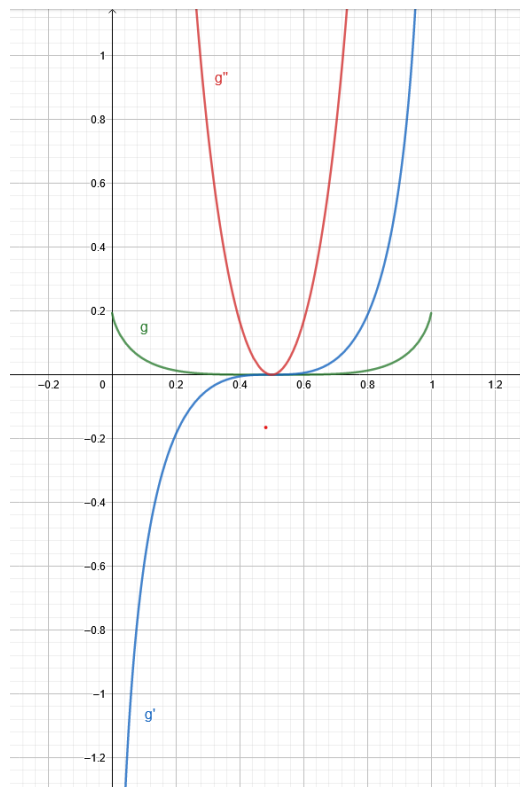
We now analyze

$$\begin{aligned}g(p) &= f\left(p, \frac{1}{2}\right) = p \cdot \log(2p) + (1 - p) \cdot \log(2 - 2p) - 2 \cdot \left(p - \frac{1}{2}\right)^2 \\ g'(p) &= \log(2p) + p \cdot \frac{1}{2p} \cdot 2 + (-1) \cdot \log(2 - 2p) + (1 - p) \cdot \frac{1}{2 - 2p} \cdot (-2) - 4 \cdot \left(p - \frac{1}{2}\right) \\ &= \log(2p) + 1 - \log(2 - 2p) - 1 - 4p + 2 \\ &= \log(2p) - \log(2 - 2p) - 4p + 2. \\ g''(p) &= \frac{1}{2p} \cdot 2 - \frac{1}{2 - 2p} \cdot (-2) - 4 \\ &= \frac{1}{p} + \frac{1}{1 - p} - 4 = \frac{1 - p + p - 4p + 4p^2}{p \cdot (1 - p)} \\ &= \frac{4p^2 - 4p + 1}{p - p^2} = \frac{4\left(p - \frac{1}{2}\right)^2}{p \cdot (1 - p)}.\end{aligned}$$

For $p \in (0, 1)$, the denominator of the second derivative is always positive and therefore the sign depends only on the nominator. As the nominator is also non-negative for $p \in (0, 1)$, the second derivative is always non-negative and the first derivative is therefore a non-decreasing function.

$$\begin{aligned}g'\left(\frac{1}{2}\right) &= \log 1 - \log 1 - 2 + 2 = 0 \\ g\left(\frac{1}{2}\right) &= \frac{1}{2} \cdot \log 1 + \frac{1}{2} \cdot \log 1 - 0 = 0 \\ g(0) &= 0 \cdot \log 0 + 1 \cdot \log 2 = \infty \\ g(1) &= 1 \cdot \log 2 + 0 \cdot \log 0 = \infty\end{aligned}$$

The figure sums up our analysis:

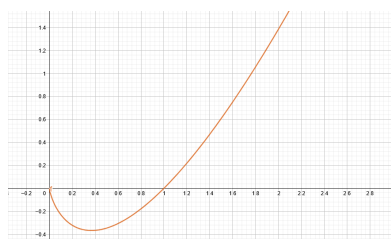


□

Therefore, the Pinsker's inequality holds for two arbitrary Bernoulli distributions. For the general case, we will need the log sum inequality and the information processing inequality:

Lemma 2.2. Log sum inequality Let $p_1, p_2, \dots, p_n, q_1, q_2, \dots, q_n \in \mathbb{R}_0^+$ be non-negative real numbers. Let $p = \sum_{i=1}^n p_i$ and $q = \sum_{i=1}^n q_i$. Then

$$\sum_{i=1}^n p_i \log \frac{p_i}{q_i} \geq p \log \frac{p}{q}.$$



Proof. Set $f(x) = x \log x$. Notice that f is a convex function. Then,

$$\begin{aligned}
\sum_{i=1}^n p_i \log \frac{p_i}{q_i} &= \sum_{i=1}^n p_i \frac{q_i}{q_i} \log \frac{p_i}{q_i} \\
&= \sum_{i=1}^n q_i f\left(\frac{p_i}{q_i}\right) \\
&= q \sum_{i=1}^n \frac{q_i}{q} f\left(\frac{p_i}{q_i}\right) \\
&\geq q \cdot f\left(\sum_{i=1}^n \frac{q_i}{q} \cdot \frac{p_i}{q_i}\right) \\
&= q \cdot f\left(\frac{1}{q} \sum_{i=1}^n p_i\right) = q \cdot f\left(\frac{p}{q}\right) \\
&= q \cdot \frac{p}{q} \cdot \log \frac{p}{q} \\
&= p \log \frac{p}{q}.
\end{aligned}$$

□

Lemma 2.3. Information processing inequality.

For any function $f : S \rightarrow S'$ and probability distributions $X : 2^S \rightarrow [0, 1]$ and $Y : 2^S \rightarrow [0, 1]$ defined over S , define

$$\begin{aligned}
X' : 2^{S'} &\rightarrow [0, 1], \\
Y' : 2^{S'} &\rightarrow [0, 1].
\end{aligned}$$

For every $i \in S'$, define

$$\begin{aligned}
X'(i) &= X(f^{-1}(i)) = \sum_{w \in f^{-1}(i)} X(w), \\
Y'(i) &= Y(f^{-1}(i)) = \sum_{w \in f^{-1}(i)} Y(w).
\end{aligned}$$

If X' and Y' are probability distributions, then

$$D_{\text{KL}}(X' || Y') \leq D_{\text{KL}}(X || Y).$$

Proof.

$$\begin{aligned}
D_{\text{KL}}(X || Y) &= \sum_{w \in S} X(w) \log \frac{X(w)}{Y(w)} \\
&= \sum_{i \in S'} \sum_{w \in f^{-1}(i)} X(w) \log \frac{X(w)}{Y(w)} \\
&\geq \sum_{i \in S'} X'(i) \log \frac{X'(i)}{Y'(i)} \\
&= D_{\text{KL}}(X' || Y').
\end{aligned}$$

□

Proof. Pinsker's inequality. Given probability distributions $P(x)$ and $Q(x)$ from discrete probability spaces defined over the same S , define $f : S \rightarrow \{0, 1\}$

$$f(w) = \begin{cases} 1 & P(w) \leq Q(w), \\ 0 & P(w) > Q(w). \end{cases}$$

Define probability distributions $P', Q' : 2^{\{0,1\}} \rightarrow [0, 1]$ for $i \in \{0, 1\}$ as

$$\begin{aligned} P'(i) &= P(f^{-1}(i)) = \sum_{w \in f^{-1}(i)} P(w), \\ Q'(i) &= Q(f^{-1}(i)) = \sum_{w \in f^{-1}(i)} Q(w), \\ P'(0) &= \sum_{\{w \in S \mid P(w) > Q(w)\}} P(w), \\ Q'(0) &= \sum_{\{w \in S \mid P(w) > Q(w)\}} Q(w), \\ P'(1) &= \sum_{\{w \in S \mid P(w) \leq Q(w)\}} P(w), \\ Q'(1) &= \sum_{\{w \in S \mid P(w) \leq Q(w)\}} Q(w). \end{aligned}$$

From this follows that $P'(0) > Q'(0)$ and $P'(1) \leq Q'(1)$.

As P' and Q' are Bernoulli distributions, we know that $D_{\text{KL}}(P' \| Q') \geq \frac{1}{2} \|P' - Q'\|_1^2$ by Pinsker's inequality for Bernoulli distributions.

Also,

$$\begin{aligned} \|P - Q\|_1 &= \sum_{w \in S} |P(w) - Q(w)| \\ &= \sum_{w \in f^{-1}(0)} (P(w) - Q(w)) + \sum_{w \in f^{-1}(1)} (Q(w) - P(w)) \\ &= P'(0) - Q'(0) + Q'(1) - P'(1) \\ &= |P'(0) - Q'(0)| + |P'(1) - Q'(1)| \\ &= \|P' - Q'\|_1. \end{aligned}$$

Therefore, $D_{\text{KL}}(P' \| Q') \geq \frac{1}{2} \|P - Q\|_1^2$.

By information processing inequality, we know that

$$D_{\text{KL}}(P' \| Q') \leq D_{\text{KL}}(P \| Q).$$

And that is all, folks!

$$D_{\text{KL}}(P \| Q) \geq D_{\text{KL}}(P' \| Q') \geq \frac{1}{2} \|P' - Q'\|_1^2 = \frac{1}{2} \|P - Q\|_1^2.$$

□

3 Upper Bound

There does not exist such a nice lower bound for KL divergence for a simple reason.

3.1 A counterexample

Theorem 3.1. Kullback-Leibler divergence is not upper bounded by the L_1 metric. Formally, for every $\varepsilon > 0$, there exist probability distributions P_ε and Q such that:

$$\|P - Q\|_1 \leq \varepsilon, \text{ but } D_{\text{KL}}(P \| Q) = \infty.$$

Proof. Define $P(x)$ and Q as

$$\begin{aligned} S &= \{a, b\} \\ Q(a) &= 0, Q(b) = 1 \\ P_\varepsilon(a) &= \frac{\varepsilon}{2}, P_\varepsilon(b) = 1 - \frac{\varepsilon}{2} \end{aligned}$$

Then,

$$\begin{aligned} \|P - Q\|_1 &= \varepsilon, \\ D_{\text{KL}}(P||Q) &= \frac{\varepsilon}{2} \cdot \log \frac{\frac{\varepsilon}{2}}{0} = \infty. \end{aligned}$$

□

3.2 A proof

Theorem 3.2. For two probability distributions $P(x)$ and $Q(x)$ that are defined over the same S , it holds that

$$D_{\text{KL}}(P||Q) \leq \frac{1}{2\alpha_Q} \|P - Q\|_1^2,$$

where

$$\alpha_Q = \min_{x \in S} Q(x).$$

4 Misc

4.1 Total variation distance vs L_1 norm

Theorem 4.1. Scheffé's lemma.

For two probability distributions $P(x)$ and $Q(x)$ that are defined over the same S , it holds that

$$\delta(P, Q) = \frac{1}{2} \|P - Q\|_1.$$

Proof. Let's refresh the definition of the total variation distance:

$$\delta(P, Q) = \max_{A \in 2^S} |P(A) - Q(A)|. \quad (5)$$

Denote by $G = \{x \in S \mid P(x) \geq Q(x)\}$. Try to find $A \subset S$ such that $P(A) - Q(A)$ is maximized. Intuitively, it is the case when $A = G$.

Now, try to find $A' \subset S$ such that $Q(A) - P(A)$ is maximized. Intuitively, it is the case when $A' = S \setminus G$.

Therefore, the subset A in $\delta(P, Q) = \max_{A \in 2^S} |P(A) - Q(A)|$ is either $A = G$ or $A' = S \setminus G$. We will show that the maximum is obtained at both A and A' :

$$P(G) - Q(G) = (1 - P(S \setminus G)) - (1 - Q(S \setminus G)) = Q(S \setminus G) - P(S \setminus G).$$

So if A maximizes $P(X) - Q(X)$, A' maximizes $Q(X) - P(X)$ and if A' maximizes $P(X) - Q(X)$, then A maximizes $Q(X) - P(X)$.

Now,

$$\begin{aligned} \|P - Q\|_1 &= \sum_{x \in S} |P(x) - Q(x)| \\ &= \sum_{x \in G} (P(x) - Q(x)) + \sum_{x \in S \setminus G} (Q(x) - P(x)) \\ &= \delta(P, Q) + \delta(P, Q) \\ &= 2 \cdot \delta(P, Q) \end{aligned}$$

□

4.2 Inverse Pinsker inequality

Let P and Q be probability distributions on the finite set A . Let $A_+ = \{a : Q(a) > 0\}$ and let $\alpha_Q = \min_{a \in A_+} Q(a)$.

How to prove that if $D(P||Q) < \infty$ then

$$D(P||Q) \leq \frac{d^2(P, Q)}{\alpha_Q \cdot \ln 2},$$

where $d(P, Q)$ is the variational distance of distributions P and Q , i.e., $d(P, Q) = \sum_{a \in A} |P(a) - Q(a)|$.

I was given a hint that first should prove that:

$$D(P||Q) \leq \sum_{a \in A_+} \frac{P(a)}{\ln 2} \left(\frac{P(a)}{Q(a)} - 1 \right) = \frac{1}{\ln 2} \sum_{a \in A_+} \frac{|P(a) - Q(a)|^2}{Q(a)}.$$

statistics information-theory

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asked Mar 1, 2020 at 15:25

 tefodos139
1

What did you try, and where are you stuck? Can you prove the hinted inequality? Can you conclude the argument assuming the hint? - stochasticboy321 Mar 5, 2020 at 4:07

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For the Hint from the right hand side to the middle is easy, hint distract the squared expression.

$$\begin{aligned} \sum_{a \in A_+} \frac{P(a)}{\ln 2} \left(\frac{P(a)}{Q(a)} - 1 \right) &= \frac{1}{\ln 2} \sum_{a \in A_+} P(a) \exp\left(\ln \frac{P(a)}{Q(a)}\right) - 1 \geq \frac{1}{\ln 2} \exp\left(\sum_{a \in A_+} P(a) \ln\left(\frac{P(a)}{Q(a)}\right)\right) - 1 \\ &\geq \exp(D(P||Q)) - 1 \\ &\geq 1 + D(P||Q) - 1 = D(P||Q) \end{aligned}$$

For the reversed pinker's;

$$\begin{aligned} D(P||Q) &\leq \frac{1}{\ln 2} \sum_{a \in A_+} \frac{|P(a) - Q(a)|^2}{Q(a)} \\ &\leq \frac{1}{\ln 2} \sum_{a \in A_+} \frac{|P(a) - Q(a)|^2}{\min_{a \in A_+} Q(a)} \leq \frac{\max_{a \in A_+} |P(a) - Q(a)| \cdot \sum_{a \in A_+} |P(a) - Q(a)|}{\alpha_Q \cdot \ln 2} \leq \frac{d^2(P, Q)}{\alpha_Q \cdot \ln 2} \end{aligned}$$

The last inequality you can deduce it after using scheffe's theorem for variational distance.