Unconstrained Optimization Overview

Notation

In what follows, we will work with vectors in \mathbb{R}^n .

The vectors will be (usually) denoted by $x \in \mathbb{R}^n$.

We often consider sequences of vectors, $x_0, x_1, \ldots, x_k, \ldots$

The index k will usually indicate that x_k is the k-the vector in a sequence.

When we talk (relatively rarely) about components of vectors, we use *i* as an index, i.e., x_i will be the *i*-th component of $x \in \mathbb{R}^n$.

We denote by ||x|| the Euclidean norm of x.

We denote by $||x||_{\infty}$ the \mathcal{L}^{∞} norm giving the maximum of absolute values of components of x.

We ocasionally use the matrix morn ||A||, consistent with the Euclidean norm, defined by

$$||A|| = \sup_{||x||=1} ||Ax|| = \sqrt{\lambda_1}$$

Here λ_1 is the largest eigenvalue of $A^{\top}A$.

How to Recognize (Local) Minimum

How do we verify that $x^* \in \mathbb{R}^n$ is a minimizer of f?



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Technically, we should examine *all* points in the immediate vicinity if one has a smaller value (impractical).

Assuming the smoothness of f, we may benefit from the "stable" behavior of f around x^* .

Derivatives and Gradients

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The gradient of $f : \mathbb{R}^n \to \mathbb{R}$, denoted by $\nabla f(x)$, is a column vector of first-order partial derivatives of the function concerning each variable:

$$abla f(x) = \left[rac{\partial f}{\partial x_1}, rac{\partial f}{\partial x_2}, \dots, rac{\partial f}{\partial x_n}
ight]^+,$$

Where each partial derivative is defined as the following limit:

$$\frac{\partial f}{\partial \mathbf{x}_{i}} = \lim_{\varepsilon \to 0} \frac{f(x_{1}, \dots, x_{i} + \varepsilon, \dots, x_{n}) - f(x_{1}, \dots, x_{i}, \dots, x_{n})}{\varepsilon}$$

Gradient



The gradient is a vector pointing in the direction of the most significant function increase from the current point.

Gradient

Consider the following function of two variables:

$$f(x_1, x_2) = x_1^3 + 2x_1x_2^2 - x_2^3 - 20x_1.$$

$$\nabla f(x_1, x_2) = \begin{bmatrix} 3x_1^2 + 2x_2^2 - 20 \\ 4x_1x_2 - 3x_2^2 \end{bmatrix}$$



Directional Derivatives vs Gradient

The rate of change in a direction p is quantified by a directional derivative, defined as

$$abla_{p}f(x) = \lim_{\varepsilon \to 0} \frac{f(x + \varepsilon p) - f(x)}{\varepsilon}$$

We can find this derivative by projecting the gradient onto the desired direction p using the dot product $\nabla_p f(x) = (\nabla f(x))^\top p$



(Here, we assume continuous partial derivatives.)

Geometry of Gradient

Consider the geometric interpretation of the dot product:

 $\nabla_{p}f(x) = (\nabla f(x))^{\top}p = ||\nabla f|| \, ||p|| \cos \theta$

Here θ is the angle between ∇f and p.

Geometry of Gradient

Consider the geometric interpretation of the dot product:

$$\nabla_{p}f(x) = (\nabla f(x))^{\top}p = ||\nabla f|| \, ||p|| \cos \theta$$

Here θ is the angle between ∇f and p.

The directional derivative is maximized by $\theta = 0$, i.e. when ∇f and p point in the same direction.



Hessian

Taking derivative twice, possibly w.r.t. different variables, gives the Hessian of \boldsymbol{f}

$$\nabla^2 f(x) = H(x) = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_1 \partial x_n} \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2^2} & \cdots & \frac{\partial^2 f}{\partial x_2 \partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1} & \frac{\partial^2 f}{\partial x_n \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_n^2} \end{bmatrix}$$

Note that the Hessian is a function which takes $x \in \mathbb{R}^n$ and gives a $n \times n$ -matrix of second derivatives of f.

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We have

$$H_{ij}=\frac{\partial^2 f}{\partial x_i\partial x_j}.$$

If f has continuous second partial derivatives, then H is symmetric, i.e., $H_{ij} = H_{ji}$.

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Let x be fixed and let g(t) = f(x + tp) and let $h_i(t) = \frac{\partial f}{\partial x_i}(x + tp)$ for $t \in \mathbb{R}$.

What exactly are g'(0) and g''(0)?

$$g'(t) = f(x+tp)' = [\nabla f(x+tp)]^{\top}p = \sum_{i=1}^{n} h_i(t)p_i$$

$$h'_{i}(t) = \left[\nabla \frac{\partial f}{\partial x_{i}}(x+tp)\right]^{\top} p = \sum_{j=1}^{n} \left(\frac{\partial f}{\partial x_{i}\partial x_{j}}(x+tp)\right) p_{j}$$
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$$g''(t) = \sum_{i=1}^{n} h'_i(t) p_i = \sum_{i=1}^{n} [H(x+tp)p]_i p_i = p^{\top} H(x+tp)p$$

Thus,

$$g''(0) = p^\top H(x)p.$$

Principal Curvature Directions

Fix x and consider H = H(x). Consider unit eigenvectors \hat{v}_k of H:

 $H\hat{v}_k = \kappa_k\hat{v}_k$

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Principal Curvature Directions

Fix x and consider H = H(x). Consider unit eigenvectors \hat{v}_k of H:

$$H\hat{v}_k = \kappa_k \hat{v}_k$$

For symmetric H, the unit eigenvectors form an orthonormal basis, and there is a rotation matrix R such that

$$H = RDR^{-1} = RDR^{\top}$$

Here *D* is diagonal with $\kappa_1, \ldots, \kappa_n$ on the diagonal.

If $\kappa_1 \geq \cdots \geq \kappa_n$, the direction of \hat{v}_1 is the maximum curvature direction of f at x.



Consider $f(x) = x^{\top} H x$ where

$$H = \begin{pmatrix} 4/3 & 0 \\ 0 & 1 \end{pmatrix}$$

The eigenvalues are

 $\kappa_1 = 4/3$ $\kappa_2 = 1$

Their corresponding eigenvectors are $(1,0)^{\top}$ and $(0,1)^{\top}$.



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Note that

$$f(x) = \kappa_1 x_1^2 + \kappa_2 x_2^2$$

Considering a direction vector p we get

$$g(t) = f(0 + tp) = t^2 (\kappa_1 p_1^2 + \kappa_2 p_2^2)$$

which is a parabola with $g'' = 2 \left(\kappa_1 p_1^2 + \kappa_2 p_2^2 \right)$.



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$$\kappa_1 = \frac{1}{6}(7 + \sqrt{5})$$
 $\kappa_2 = \frac{1}{6}(7 - \sqrt{5})$

Their corresponding eigenvectors are

$$\hat{\mathbf{v}}_{1} = \left(rac{1}{2}(1+\sqrt{5}),1
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Here $(\hat{v}_1 \ \hat{v}_2)$ is a 2 × 2 matrix whose columns are \hat{v}_1, \hat{v}_2 .



Hessian Visualization Example

Consider

$$f(x_1, x_2) = x_1^3 + 2x_1x_2^2 - x_2^3 - 20x_1.$$

And it's Hessian.

$$H(x_1, x_2) = \begin{bmatrix} 6x_1 & 4x_2 \\ 4x_2 & 4x_1 - 6x_2 \end{bmatrix}.$$



Theorem 1 (Taylor)

Suppose that $f : \mathbb{R}^n \to \mathbb{R}$ is twice continuously differentiable and that $p \in \mathbb{R}^n$. Then, we have

$$f(x+p) = f(x) + \nabla f(x)^{T} p + \frac{1}{2} p^{T} H(x) p + o(||p||^{2}).$$

Here $H = \nabla^2 f$ is the Hessian of f.

First-Order Necessary Conditions

Theorem 2

If x^* is a local minimizer and f is continuously differentiable in an open neighborhood of x^* , then $\nabla f(x^*) = 0$.





Note that $\nabla f(x^*) = 0$ does not tell us whether x^* is a minimizer, maximizer, or a saddle point.

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All comes down to the *definiteness* of $H := H(x^*)$.

- H is positive definite if p^THp > 0 for all p iff all eigenvalues of H are positive
- H is positive semi-definite if p[⊤]Hp ≥ 0 for all p iff all eigenvalues of H are nonnegative
- H is negative semi-definite if p[⊤]Hp ≤ 0 for all p iff all eigenvalues of H are nonpositive
- ► H is negative definite if p^T Hp < 0 for all p iff all eigenvalues of H are negative
- H is indefinite if it is not definite in the above sense iff H has at least one positive and one negative eigenvalue.

Definiteness



Second-Order Necessary Condition

Theorem 3 (Second-Order Necessary Conditions)

If x^* is a local minimizer of f and $\nabla^2 f$ is continuous in a neighborhood of x^* , then $\nabla f(x^*) = 0$ and $\nabla^2 f(x^*)$ is positive semidefinite.

Theorem 4 (Second-Order Sufficient Conditions)

Suppose that $\nabla^2 f$ is continuous in a neighborhood of x^* and that $\nabla f(x^*) = 0$ and $\nabla^2 f(x^*)$ is positive definite. Then x^* is a strict local minimizer of f.



Example

Consider the following function of two variables:

$$f(x_1, x_2) = 0.5x_1^4 + 2x_1^3 + 1.5x_1^2 + x_2^2 - 2x_1x_2.$$

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Consider the gradient equal to zero:

$$\nabla f = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \end{bmatrix} = \begin{bmatrix} 2x_1^3 + 6x_1^2 + 3x_1 - 2x_2 \\ 2x_2 - 2x_1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

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The solution of this equation yields three points:

$$x_{\mathcal{A}} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \quad x_{\mathcal{B}} = \begin{bmatrix} -\frac{3}{2} - \frac{\sqrt{7}}{2} \\ -\frac{3}{2} - \frac{\sqrt{7}}{2} \end{bmatrix}, \quad x_{\mathcal{C}} = \begin{bmatrix} \frac{\sqrt{7}}{2} - \frac{3}{2} \\ \frac{\sqrt{7}}{2} - \frac{3}{2} \end{bmatrix}$$

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To classify x_A, x_B, x_C , we need to compute the Hessian matrix:

$$H(x_1, x_2) = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2^2} \end{bmatrix} = \begin{bmatrix} 6x_1^2 + 12x_1 + 3 & -2 \\ -2 & 2 \end{bmatrix}$$

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The Hessian, at the first point, is

$$H(x_{\mathcal{A}}) = \begin{bmatrix} 3 & -2 \\ -2 & 2 \end{bmatrix},$$

whose eigenvalues are $\kappa_1 \approx 0.438$ and $\kappa_2 \approx 4.561$. Because both eigenvalues are positive, this point is a local minimum.

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For the second point,

$$H(x_B) = \left[egin{array}{cc} 3(3+\sqrt{7}) & -2 \ -2 & 2 \end{array}
ight].$$

The eigenvalues are $\kappa_1 \approx 1.737$ and $\kappa_2 \approx 17.200$, so this point is another local minimum.

Consider the following function of two variables:

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To classify x_A, x_B, x_C , we need to compute the Hessian matrix:

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For the third point,

$$H(x_C) = \left[\begin{array}{rrr} 9 - 3\sqrt{7} & -2\\ -2 & 2 \end{array}\right]$$

The eigenvalues for this Hessian are $\kappa_1 \approx -0.523$ and $\kappa_2 \approx 3.586$, so this point is a saddle point.



Proofs of Some Theorems Optional

Taylor's Theorem

To prove the theorems characterizing minima/maxima, we need the following form of Taylor's theorem:

Theorem 5 (Taylor)

Suppose that $f : \mathbb{R}^n \to \mathbb{R}$ is continuously differentiable and that $p \in \mathbb{R}^n$. Then we have that.

$$f(x+p) = f(x) + \nabla f(x+tp)^T p,$$

for some $t \in (0, 1)$. Moreover, if f is twice continuously differentiable, we have that

$$f(x+p) = f(x) + \nabla f(x)^T p + \frac{1}{2} p^T \nabla^2 f(x+tp)p,$$

for some $t \in (0, 1)$.

Proof of Theorem 2 (Optional)

We prove that if x^* is a local minimizer and f is continuously differentiable in an open neighborhood of x^* , then $\nabla f(x^*) = 0$.

Suppose for contradiction that $\nabla f(x^*) \neq 0$. Define the vector $p = -\nabla f(x^*)$ and note that $p^T \nabla f(x^*) = - \|\nabla f(x^*)\|^2 < 0$. Because ∇f is continuous near x^* , there is a scalar T > 0 such that

$$p^T
abla f(x^* + tp) < 0,$$
 for all $t \in [0, T]$

For any $ar{t} \in (0, T]$, we have by Taylor's theorem that

$$f\left(x^{*}+ar{t}p
ight)=f\left(x^{*}
ight)+ar{t}p^{T}
abla f\left(x^{*}+tp
ight), \hspace{0.5cm} ext{ for some }t\in(0,ar{t}).$$

Therefore, $f(x^* + \bar{t}p) < f(x^*)$ for all $\bar{t} \in (0, T]$. We have found a direction leading away from x^* along which f decreases, so x^* is not a local minimizer, and we have a contradiction.

Proof of Theorem 3 (Optional)

We prove that if x^* is a local minimizer of f and $\nabla^2 f$ is continuous in an open neighborhood of x^* , then $\nabla f(x^*) = 0$ and $\nabla^2 f(x^*)$ is positive semidefinite.

We know that $\nabla f(x^*) = 0$. For contradiction, assume that $\nabla^2 f(x^*)$ is not positive semidefinite.

Then we can choose a vector p such that $p^T \nabla^2 f(x^*) p < 0$.

As $\nabla^2 f$ is continuous near x^* , $p^T \nabla^2 f(x^* + tp) p < 0$ for all $t \in [0, T]$ where T > 0.

By Taylor we have for all $\overline{t} \in (0, T]$ and some $t \in (0, \overline{t})$

$$f(x^* + \bar{t}p) = f(x^*) + \bar{t}p^T \nabla f(x^*) + \frac{1}{2} \bar{t}^2 p^T \nabla^2 f(x^* + tp) p < f(x^*).$$

Thus, x^* is not a local minimizer.

Proof of Theorem 4 (Optional)

We prove the following: Suppose that $\nabla^2 f$ is continuous in an open neighborhood of x^* and that $\nabla f(x^*) = 0$ and $\nabla^2 f(x^*)$ is positive definite. Then x^* is a strict local minimizer of f.

Because the Hessian is continuous and positive definite at x^* , we can choose a radius r > 0 so that $\nabla^2 f(x)$ remains positive definite for all x in the open ball $\mathcal{D} = \{z \mid ||z - x^*|| < r\}$. Taking any nonzero vector p with ||p|| < r, we have $x^* + p \in \mathcal{D}$ and so

$$f(x^* + p) = f(x^*) + p^T \nabla f(x^*) + \frac{1}{2} p^T \nabla^2 f(z) p$$

= $f(x^*) + \frac{1}{2} p^T \nabla^2 f(z) p$,

where $z = x^* + tp$ for some $t \in (0, 1)$. Since $z \in D$, we have $p^T \nabla^2 f(z) p > 0$, and therefore $f(x^* + p) > f(x^*)$, giving the result.

Unconstrained Optimization Algorithms

Search Algorithms

We consider algorithms that

- Start with an initial guess x₀
- ▶ Generate a sequence of points *x*₀, *x*₁,...
- Stop when no progress can be made or when a minimizer seems approximated with sufficient accuracy.

To compute x_{k+1} the algorithms use the information about f at the previous iterates x_0, x_1, \ldots, x_k .

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There are two overall strategies:

- Line search
- Trust region

Line Search Overview

To compute x_{k+1} , a line search algorithm chooses

- \blacktriangleright direction p_k
- step size α_k

and computes

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direction p_k

• step size α_k

and computes

 $x_{k+1} = x_k + \alpha_k p_k$

The vector p_k should be a *descent* direction, i.e., a direction in which f decreases locally.

 α_k is selected to approximately solve

$$\min_{\alpha>0}f(x_k+\alpha p_k)$$

However, typically, an exact solution is expensive and unnecessary. Instead, line search algorithms inspect a limited number of trial step lengths and find one that decreases f appropriately (see later).



A descent direction does not have to be followed to the minimum.



Trust Region

To compute x_{k+1} , a trust region algorithm chooses

• model function m_k whose behavior near x_k is similar to f

▶ a trust region $R \subseteq \mathbb{R}^n$ around x_k . Usually R is the ball defined by $||x - x_k|| \leq \Delta$ where $\Delta > 0$ is trust region radius.

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The model m_k is usually derived from the Taylor's theorem.

$$m_k(x_k+p) = f_k + p^T \nabla f_k + \frac{1}{2} p^T B_k p$$

Where B_k approximates the Hessian of f at x_k .



Line Search Methods

Line Search

For setting the step size, we consider

- Armijo condition and backtracking algorithm
- strong Wolfe conditions and bracketing & zooming

Line Search

For setting the step size, we consider

- Armijo condition and backtracking algorithm
- strong Wolfe conditions and bracketing & zooming

For setting the direction, we consider

- Gradient descent
- Newton's method
- quasi-Newton methods (BFGS)
- (Conjugate gradients)

We start with the step size.

Step Size

Assume

 $x_{k+1} = x_k + \alpha_k p_k$

Where p_k is a descent direction

 $p_k^\top \nabla f_k < 0$



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$$\phi(\alpha) = f(x_k + \alpha p_k)$$



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Define

$$\phi(\alpha) = f(x_k + \alpha p_k)$$

We know that

 $\phi'(\alpha) = \nabla f(x_k + \alpha p_k)^\top p_k$ which means $\phi'(0) = \nabla f_k^\top p_k$

Note that $\phi'(0)$ must be negative as p_k is a descent direction.

Armijo Condition

The sufficient decrease condition (aka Armijo condition)

$$\phi(\alpha) \le \phi(\mathbf{0}) + \alpha \left(\mu_1 \phi'(\mathbf{0})\right)$$

where μ_1 is a constant such that $0 < \mu_1 \leq 1$



In practice, μ_1 is several orders smaller than 1, typically $\mu_1 = 10^{-4}$.

Backtracking Line Search Algorithm

Algorithm 1 Backtracking Line Search Input: $\alpha_{init} > 0, 0 < \mu_1 < 1, 0 < \rho < 1$ Output: α^* satisfying sufficient decrease condition 1: $\alpha \leftarrow \alpha_{init}$ 2: while $\phi(\alpha) > \phi(0) + \alpha \mu_1 \phi'(0)$ do 3: $\alpha \leftarrow \rho \alpha$

4: end while

The parameter ρ is typically set to 0.5. It can also be a variable set by a more sophisticated method (interpolation).

The α_{init} depends on the method for setting the descent direction p_k . For Newton and quasi-Newton, it is 1.0, but for other methods, it might be different.

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- The guess for the initial step immediately satisfies sufficient decrease. However, the function's slope is still highly negative, and we could have decreased the function value by much more if we had taken a more significant step. In this case, our guess for the initial step is far too small.

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- The guess for the initial step immediately satisfies sufficient decrease. However, the function's slope is still highly negative, and we could have decreased the function value by much more if we had taken a more significant step. In this case, our guess for the initial step is far too small.

Even if our original step size is not too far from an acceptable one, the basic backtracking algorithm ignores any information we have about the function values and gradients. It blindly takes a reduced step based on a preselected ratio ρ .
Backtracking Example

$$f(x_1, x_2) = 0.1x_1^6 - 1.5x_1^4 + 5x_1^2 + 0.1x_2^4 + 3x_2^2 - 9x_2 + 0.5x_1x_2$$

2.5 2 1.5 1 0.5 -3 -2 -1 0 1 2 3 x_1 x_1

 x_2

 $\mu_1 = 10^{-4}$ and $\rho = 0.7$.



We want to prevent too short of steps and to "motivate" the search to move closer to the minimum.

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We introduce the sufficient curvature condition

 $\left|\phi'(\alpha)\right| \le \mu_2 \left|\phi'(0)\right|$

where $\mu_1 < \mu_2 < 1$ is a constant.



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Typical values of μ_2 range from 0.1 to 0.9, depending on the direction setting method.

As μ_2 tends to 0, the condition enforces $\phi'(\alpha) = 0$, which would yield an exact line search.

Strong Wolfe Conditions

Putting together Armijo and sufficient curvature conditions, we obtain *strong Wolfe conditions*

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Sufficient curvature condition $\left|\phi'(\alpha)\right| \leq \mu_2 \left|\phi'(0)\right|$ $\phi'(0$ $\phi(0)$ $\mu_1 \phi'(0)$ Sufficient decrease line $\phi(\alpha)$ $\mu_2 \phi'(0)$ $\alpha = 0$ α Acceptable range Acceptable range

Satisfiability of Strong Wolfe Conditions

Theorem 6

Suppose $f : \mathbb{R}^n \to \mathbb{R}$ is continuously differentiable. Let p_k be a descent direction at x_k , and assume that f is bounded below along the ray $\{x_k + \alpha p_k \mid \alpha > 0\}$. Then, if $0 < \mu_1 < \mu_2 < 1$, step length intervals exist that satisfy the strong Wolfe conditions.



Convergence of Line Search

Denote by θ_k the angle between p_k and $-\nabla f_k$, i.e., satisfying

$$\cos \theta_k = \frac{-\nabla f_k^T p_k}{\|\nabla f_k\| \, \|p_k\|}$$

Convergence of Line Search

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$$\cos \theta_k = \frac{-\nabla f_k^T \rho_k}{\|\nabla f_k\| \, \|\rho_k\|}$$

Recall that f is L-smooth for some L > 0 if

$$\|
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Theorem 7 (Zoutendijk)

Consider $x_{k+1} = x_k + \alpha_k p_k$, where p_k is a descent direction and α_k satisfies the strong Wolfe conditions. Suppose that f is bounded below, continuously differentiable, and L-smooth. Then

$$\sum_{k\geq 0}\cos^2\theta_k \|\nabla f_k\|^2 < \infty.$$

Line Search Algorithm

How can we find a step size that satisfies strong Wolfe conditions?

How can we find a step size that satisfies *strong* Wolfe conditions? Use a bracketing and zoom algorithm, which proceeds in the following two phases:

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Line Search Algorithm

How can we find a step size that satisfies strong Wolfe conditions?

Use a bracketing and zoom algorithm, which proceeds in the following two phases:

- 1. The bracketing phase finds an interval within which we are certain to find a point that satisfies the strong Wolfe conditions.
- The zooming phase finds a point that satisfies the strong Wolfe conditions within the interval provided by the bracketing phase.

Algorithm 2 Bracketing

Input: $\alpha_1 > 0$ and α_{max} 1: Set $\alpha_0 \leftarrow 0$ 2: $i \leftarrow 1$ 3: repeat Evaluate $\phi(\alpha_i)$ 4: if $\phi(\alpha_i) > \phi(0) + \alpha_i \mu_1 \phi'(0)$ or $[\phi(\alpha_i) \ge \phi(\alpha_{i-1})$ and i > 15: then $\alpha^* \leftarrow \mathbf{zoom}(\alpha_{i-1}, \alpha_i)$ and stop 6: end if 7: Evaluate $\phi'(\alpha_i)$ 8: if $|\phi'(\alpha_i)| < \mu_2 |\phi'(0)|$ then 9: set $\alpha^* \leftarrow \alpha_i$ and stop 10: else if $\phi'(\alpha_i) > 0$ then 11: set $\alpha^* \leftarrow \mathbf{zoom}(\alpha_i, \alpha_{i-1})$ and stop 12: end if 13: Choose $\alpha_{i+1} \in (\alpha_i, \alpha_{\max})$ 14: $i \leftarrow i + 1$ 15: 16: **until** a condition is met

Explanation of Bracketing

Note that the sequence of trial steps α_i is monotonically increasing.

Note that the sequence of trial steps α_i is monotonically increasing.

Note that **zoom** is called when one of the following conditions is satisfied:

- α_i violates the sufficient decrease condition (lines 5 and 6)
- $\phi(\alpha_i) \ge \phi(\alpha_{i-1})$ (also lines 5 and 6)
- $\phi'(\alpha_i) \ge 0$ (lines 11 and 12)

The last step increases the α_i . May use, e.g., a constant multiple.

The following algorithm keeps two step lengths: $\alpha_{\rm lo}$ and $\alpha_{\rm hi}$

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The following invariants are being preserved:

The interval bounded by α_{lo} and α_{hi} always contains one or more intervals satisfying the strong Wolfe conditions. Note that we *do not* assume α_{lo} ≤ α_{hi}

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1: function $ZOOM(\alpha_{lo}, \alpha_{hi})$

2: repeat

- 3: Set α between α_{lo} and α_{hi} using interpolation (bisection, quadratic, etc.)
- 4: Evaluate $\phi(\alpha)$
- 5: if $\phi(\alpha) > \phi(0) + \alpha \mu_1 \phi'(0)$ or $\phi(\alpha) \ge \phi(\alpha_{\mathsf{lo}})$ then
- 6: $\alpha_{hi} \leftarrow \alpha$

else

7:

- 8: Evaluate $\phi'(\alpha)$
- 9: **if** $|\phi'(\alpha)| \le \mu_2 |\phi'(0)|$ then 10: Set $\alpha^* \leftarrow \alpha$ and stop
- 10: Set $\alpha^* \leftarrow$ 11: **end if**
- 12: **if** $\phi'(\alpha)(\alpha_{hi} \alpha_{lo}) \ge 0$ then

13: $\alpha_{hi} \leftarrow \alpha_{lo}$

- 14: **end if**
- 15: $\alpha_{\mathsf{lo}} \leftarrow \alpha$
- 16: end if
- 17: **until** a condition is met
- 18: end function

Bracketing & Zooming Example

We use quadratic interpolation; the bracketing chooses $\alpha_{i+1} = 2\alpha_i$, and the sufficient curvature factor is $\mu_2 = 0.9$.



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Bracketing is achieved in the first iteration by using a significant initial step of $\alpha_{init} = 1.2$ (left). Then, zooming finds an improved point through interpolation.

Bracketing & Zooming Example

We use quadratic interpolation; the bracketing chooses $\alpha_{i+1} = 2\alpha_i$, and the sufficient curvature factor is $\mu_2 = 0.9$.



Bracketing is achieved in the first iteration by using a significant initial step of $\alpha_{init} = 1.2$ (left). Then, zooming finds an improved point through interpolation.

The small initial step of $\alpha_{init} = 0.05$ (right) does not satisfy the strong Wolfe conditions, and the bracketing phase moves forward toward a flatter part of the function.

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- A problem that can arise in the implementation is that as the optimization algorithm approaches the solution, two consecutive function values f (x_k) and f (x_{k-1}) may be indistinguishable in finite-precision arithmetic.
- Some procedures also stop if the relative change in x is close to machine accuracy or some user-specified threshold.
- The presented algorithm is implemented in https://docs.scipy.org/doc/scipy/reference/ generated/scipy.optimize.line_search.html

Unconstrained Optimization Algorithms

Descent Direction

First-Order Methods

Gradient Descent

Consider the *gradient descent* (aka *gradient descent*) method where

$$x_{k+1} = x_k + \alpha_k p_k$$
 $p_k = -\nabla f(x_k)$



Gradient Descent

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 $p_k = -\nabla f(x_k)$



Unfortunately, the gradient does not possess much information about the step size.

So usually, a normalized gradient is used to obtain the direction, and then a line search is performed:

$$x_{k+1} = x_k + \alpha_k p_k$$
 $p_k = -\frac{\nabla f(x_k)}{||\nabla f(x_k)||}$

The line search is *exact* if α_k minimizes $f(x_k + \alpha_k p_k)$. Not practical, we usually find α_k satisfying the strong Wolfe conditions.

Gradient Descent Algorithm with Line Search

Algorithm 3 Gradient Descent with Line Search

Input: x_0 starting point, $\varepsilon > 0$

Output: x^* approximation to a stationary point

1:
$$k \leftarrow 0$$

2: while
$$\|\nabla f\|_{\infty} > \varepsilon$$
 do

3:
$$p_k \leftarrow -\frac{\nabla f(x_k)}{\|\nabla f(x_k)\|}$$

4: Set
$$\alpha_{init}$$
 for line search

5:
$$\alpha_k \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$$

$$6: \qquad x_{k+1} \leftarrow x_k + \alpha_k p_k$$

7:
$$k \leftarrow k+1$$

8: end while
Gradient Descent Algorithm with Line Search

Algorithm 4 Gradient Descent with Line Search **Input:** x_0 starting point, $\varepsilon > 0$ **Output:** x^* approximation to a stationary point 1: $k \leftarrow 0$ 2: while $\|\nabla f\|_{\infty} > \varepsilon$ do $p_k \leftarrow -\frac{\nabla f(x_k)}{\|\nabla f(x_k)\|}$ 3: 4: Set α_{init} for line search 5: $\alpha_k \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ 6: $x_{k+1} \leftarrow x_k + \alpha_k p_k$ $k \leftarrow k+1$ 7. 8: end while

Here α_{init} can be estimated from the previous step size α_{k-1} by demanding similar decrease in the objective:

$$\alpha_{\text{init}} \boldsymbol{p}_{k}^{\top} \nabla f_{k} \approx \alpha_{k-1} \boldsymbol{p}_{k-1}^{\top} \nabla f_{k-1} \quad \Rightarrow \quad \alpha_{\text{init}} = \alpha_{k-1} \frac{\boldsymbol{p}_{k-1}^{\top} \nabla f_{k-1}}{\boldsymbol{p}_{k}^{\top} \nabla f_{k}}$$

Gradient Descent Algorithm with Line Search

Algorithm 5 Gradient Descent with Line Search

Input: x_0 starting point, $\varepsilon > 0$ **Output:** x^* approximation to a stationary point

1:
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2: while
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3:
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Gradient Descent Algorithm with Line Search

Algorithm 6 Gradient Descent with Line Search

Input: x_0 starting point, $\varepsilon > 0$ **Output:** x^* approximation to a stationary point 1: $k \leftarrow 0$ 2: while $\|\nabla f\|_{\infty} > \varepsilon$ do $p_k \leftarrow -\frac{\nabla f(x_k)}{\|\nabla f(x_k)\|}$ 3: 4: Set α_{init} for line search 5: $\alpha_k \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ 6: $x_{k+1} \leftarrow x_k + \alpha_k p_k$ $k \leftarrow k+1$ 7. 8: end while

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Note that p_{k+1} and p_k are always orthogonal.

$$f(x_1, x_2) = (1 - x_1)^2 + (1 - x_2)^2 + \frac{1}{2}(2x_2 - x_1^2)^2$$

Stopping: $||\nabla f||_{\infty} \leq 10^{-6}$.



The gradient descent can be prolonged.

Global Convergence with Line Search

Recall the Zoutendijk's theorem.

Denote by θ_k the angle between p_k and $-\nabla f_k$, i.e., satisfying

$$\cos \theta_k = \frac{-\nabla f_k^T p_k}{\|\nabla f_k\| \, \|p_k\|}$$

Recall that f is L-smooth on a set \mathcal{N} for some L > 0 if

$$\|
abla f(x) -
abla f(ilde{x})\| \leq L \|x - ilde{x}\|, \quad ext{ for all } x, ilde{x} \in \mathcal{N}$$

Theorem 8 (Zoutendijk)

Consider $x_{k+1} = x_k + \alpha_k p_k$, where p_k is a descent direction and α_k satisfies the strong Wolfe conditions. Suppose that f is bounded below in \mathbb{R}^n and that f is continuously differentiable in an open set \mathcal{N} containing the level set $\{x : f(x) \leq f(x_0)\}$. Assume also that f is L-smooth on \mathcal{N} . Then

$$\sum_{k\geq 0}\cos^2\theta_k \, \|\nabla f_k\|^2 < \infty.$$

Global Convergence of Gradient Descent

Assume that each α_k satisfies strong Wolfe conditions.

Global Convergence of Gradient Descent

Assume that each α_k satisfies strong Wolfe conditions.

Note that the angle θ_k between $p_k = -\nabla f_k$ and the negative gradient $-\nabla f_k$ equals 0. Hence, $\cos \theta_k = 1$.

Global Convergence of Gradient Descent

Assume that each α_k satisfies strong Wolfe conditions.

Note that the angle θ_k between $p_k = -\nabla f_k$ and the negative gradient $-\nabla f_k$ equals 0. Hence, $\cos \theta_k = 1$.

Thus, under the assumptions of Zoutendijk's theorem, we obtain

$$\sum_{k\geq 0}\cos^2\theta_k \|\nabla f_k\|^2 = \sum_{k\geq 0} \|\nabla f_k\|^2 < \infty$$

which implies that $\lim_{k\to\infty} ||\nabla f_k|| = 0$.

Local Linear Convergence of Gradient Descent

Theorem 9

Suppose that $f : \mathbb{R}^n \to \mathbb{R}$ is twice continuously differentiable, that the line search is exact, and that the descent converges to x^* where $\nabla f(x^*) = 0$ and the Hessian matrix $\nabla^2 f(x^*)$ is positive definite. Then

$$f(x_{k+1}) - f(x^*) \leq \left(\frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1}\right)^2 \left[f(x_k) - f(x^*)\right],$$

where $\lambda_1 \leq \cdots \leq \lambda_n$ are the eigenvalues of $\nabla^2 f(x^*)$.





Here $\ell_1 = 12, \ell_2 = 8, k_1 = 1, k_2 = 10, mg = 7$

Two Spring Problem - Gradient Descent



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$.

Rosenbrock Function - Gradient Descent Rosenbrock: $f(x_1, x_2) = (1 - x_1)^2 + 100(x_2 - x_1^2)^2$



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$.

The method needs evaluation of ∇f at each x_k. If f is not differentiable at x_k, subgradients can be considered (out of the scope of this course).

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- Susceptible to scaling of variables (see the paraboloid example).

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- Slow, zig-zagging, provides insufficient information for line search initialization.
- Susceptible to scaling of variables (see the paraboloid example).
- THE basis for algorithms training neural networks a huge amount of specific adjustments are developed for working with huge numbers of variables in neural networks (trillions of weights).

Unconstrained Optimization Algorithms

Descent Direction

Second-Order Methods

Consider an objective $f : \mathbb{R}^n \to \mathbb{R}$.

Assume that f is twice differentiable.

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Assume that f is twice differentiable.

Then, by the Taylor's theorem,

$$f(x_k+s) \approx f_k + \nabla f_k^{\top} s + \frac{1}{2} s^{\top} H_k s$$

where we denote the Hessian $\nabla^2 f(x_k)$ by H_k .

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where we denote the Hessian $\nabla^2 f(x_k)$ by H_k . Define

$$q(s) = f_k +
abla f_k^\top s + rac{1}{2} s^\top H_k s$$

and minimize q w.r.t. s by setting $\nabla q(s) = 0$.

Consider an objective $f : \mathbb{R}^n \to \mathbb{R}$.

Assume that f is twice differentiable.

Then, by the Taylor's theorem,

$$f(\mathbf{x}_k + \mathbf{s}) \approx f_k + \nabla f_k^\top \mathbf{s} + \frac{1}{2} \mathbf{s}^\top H_k \mathbf{s}$$

where we denote the Hessian $\nabla^2 f(x_k)$ by H_k .

Define

$$q(s) = f_k +
abla f_k^ op s + rac{1}{2} s^ op H_k s$$

and minimize q w.r.t. s by setting $\nabla q(s) = 0$. We obtain:

$$H_k s = -\nabla f_k$$

Denote by s_k the solution, and set $x_{k+1} = x_k + s_k$.

Algorithm 7 Newton's Method

Input: x_0 starting point, $\varepsilon > 0$

Output: x^* approximation to a stationary point

1:
$$k \leftarrow 0$$

2: while
$$\|\nabla f_k\|_{\infty} > \varepsilon$$
 do

3:
$$p_k \leftarrow -H_k^{-1} \nabla f(x_k)$$

4:
$$x_{k+1} \leftarrow x_k + p_k$$

5:
$$k \leftarrow k+1$$

6: end while

Newton's Method - Example

Newton's method finds the minimum of a quadratic function in a single step.



Note that the Newton's method is scale-invariant!

$$f(x_1, x_2) = (1 - x_1)^2 + (1 - x_2)^2 + \frac{1}{2} (2x_2 - x_1^2)^2$$

Stopping: $||\nabla f||_{\infty} \leq 10^{-6}$.



k = 0

k = 2

$$f(x_1, x_2) = (1 - x_1)^2 + (1 - x_2)^2 + \frac{1}{2}(2x_2 - x_1^2)^2$$

Stopping: $||\nabla f||_{\infty} \leq 10^{-6}$.



k = 2

Convergence Issues



Also, the computation of the Hessian is costly.

Theorem 10

Assume f is defined and twice differentiable and assume that ∇f is L-smooth on \mathcal{N} .

Let x_* be a minimizer of f(x) in \mathcal{N} and assume that $\nabla^2 f(x_*)$ is positive definite.

If $||x_0 - x_*||$ is sufficiently small, then $\{x_k\}$ converges quadratically to x_* .

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However, what happens if we start far away from a minimizer?

Newton's Method with Line Search

Algorithm 8 Newton's Method with Line Search **Input:** x_0 starting point, $\varepsilon > 0$

Output: x^* approximation to a stationary point

- 1: $k \leftarrow 0$
- 2: $\alpha_{\mathsf{init}} \leftarrow 1$
- 3: while $\|\nabla f_k\|_{\infty} > \varepsilon$ do
- 4: $p_k \leftarrow -H_k^{-1} \nabla f(x_k)$
- 5: $\alpha_k \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$
- $6: \qquad x_{k+1} \leftarrow x_k + p_k$
- 7: $k \leftarrow k+1$
- 8: end while



Here $\ell_1 = 12, \ell_2 = 8, k_1 = 1, k_2 = 10, mg = 7$

Two Spring Problem - Newton's Method



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$. Compare this with 32 iterations of gradient descent. Rosenbrock Function - Newton's Method Rosenbrock: $f(x_1, x_2) = (1 - x_1)^2 + 100(x_2 - x_1^2)^2$



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$. Compare this with 10,662 iterations of gradient descent.
Global Convergence of Line Search

Denote by θ_k the angle between p_k and $-\nabla f_k$, i.e., satisfying

$$\cos \theta_k = \frac{-\nabla f_k^T p_k}{\|\nabla f_k\| \, \|p_k\|}$$

Recall that f is L-smooth for some L > 0 if

$$\|
abla f(x) -
abla f(ilde{x})\| \le L \|x - ilde{x}\|, \quad \text{ for all } x, ilde{x} \in \mathbb{R}^n$$

Theorem 11 (Zoutendijk)

Consider $x_{k+1} = x_k + \alpha_k p_k$, where p_k is a descent direction and α_k satisfies the strong Wolfe conditions. Suppose that f is bounded below, continuously differentiable, and L-smooth. Then

$$\sum_{k\geq 0}\cos^2\theta_k\,\|\nabla f_k\|^2<\infty.$$

Assume that all α_k satisfy strong Wolfe conditions.

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Thus, under the assumptions of Zoutendijk's theorem, we obtain

$$\frac{1}{M^2} \sum_{k \ge 0} \|\nabla f_k\|^2 \le \sum_{k \ge 0} \cos^2 \theta_k \, \|\nabla f_k\|^2 < \infty$$

which implies that $\lim_{k\to\infty} ||\nabla f_k|| = 0$.

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which implies that $\lim_{k\to\infty} ||\nabla f_k|| = 0$.

What if H_k is not positive definite or is (nearly) singular?

Eigenvalue Modification

Consider $H_k = \nabla^2 f(x_k)$ and consider its diagonal form:

$$H_k = QDQ^T$$

Where D contains the eigenvalues of H_k on the diagonal, i.e., $D = \text{diag}(\lambda_1, \dots, \lambda_n)$ and Q is an orthogonal matrix.

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Observe that

- H_k is not positive definite iff $\lambda_i \leq 0$ for some *i*
- ▶ $||H_k||$ grows with max{ $\lambda_1, \ldots, \lambda_n$ } going to infinity.
- ► ||H_k⁻¹|| grows with min{λ₁,...,λ_n} going to 0 (i.e., the matrix becomes close to a singular matrix)

We want to prevent all three cases, i.e., make sure that for some reasonably large $\delta > 0$ we have $\lambda_i \ge \delta$ but not too large.

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We want to prevent all three cases, i.e., make sure that for some reasonably large $\delta > 0$ we have $\lambda_i \ge \delta$ but not too large.

Two questions are in order:

- What is a reasonably large δ ?
- How to modify H_k so the minimum is large enough?

Consider an example:

$$abla f(x_k) = (1, -3, 2)$$
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If used in Newton's method, we obtain the following direction:

$$p_k = -B_k^{-1} \nabla f(x_k) = (1/10, -1, -(2 \cdot 10^8))$$

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Even though f decreases along p_k , it is far from the minimum of the quadratic approximation of f.

Note that the original Newton's direction is

 $-\text{diag}(1/10, 1/3, -1)(1, -3, 2)^{\top} = (-1/10, 1, 2)$ which is completely different.

Other strategies for eigenvalue modification can be devised.

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The criteria are rather loose. The resulting matrix B_k should be

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What is ΔH_k in our example?

Various methods for computing ΔH_k have been devised in literature. Typically, it is based on some computationally cheaper decomposition than spectral decomposition (e.g., Cholesky).

Modified Newton's Method

Algorithm 9 Newton's Method with Line Search **Input:** x_0 starting point, $\varepsilon > 0$ **Output:** x^* approximation to a stationary point 1: $k \leftarrow 0$ 2: while $\|\nabla f_k\|_{\infty} > \varepsilon$ do $H_k \leftarrow \nabla^2 f(x_k)$ 3: if H_k is **not** sufficiently positive definite **then** 4: $H_k \leftarrow H_k + \Delta H_k$ so that H_k is sufficiently pos. definite 5: end if 6: Solve $H_k p_k = -\nabla f(x_k)$ for p_k 7: Set $x_{k+1} = x_k + \alpha_k p_k$, here α_k sat. the Wolfe cond. 8: $k \leftarrow k + 1$ 9. 10: end while

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In a sense, Newton's method is an impractical "ideal" with which other methods are compared.

The efficiency issues (and the necessity of second-order derivatives) will be mitigated by using quasi-Newton methods.

Recall that Newton's method step p_k in $x_{k+1} = x_k + p_k$ comes from minimization of

$$q(p) = f_k + \nabla f_k^\top p + \frac{1}{2} p^\top H_k p$$

w.r.t. p by setting $\nabla q(p) = 0$ and solving

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So Newton's method needs the second derivative (Hessian), which is computationally hard to obtain.

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Can we find a compromise?

Quasi-Newton methods use first derivatives to approximate the Hessian H_k in Newton's method with a matrix \tilde{H}_k .

Suppose we have just obtained the new point x_{k+1} after a line search starting from x_k in the direction p_k .

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First, it should be *symmetric positive definite*. To always yield decrease direction.

Second, extrapolating from the single variable secant method, we demand

$$\tilde{H}_{k+1}(x_{k+1}-x_k)=\nabla f_{k+1}-\nabla f_k$$

This is the *secant condition*.

Secant Condition

Consider the secant condition:

$$ilde{H}_{k+1}(x_{k+1}-x_k) =
abla f_{k+1} -
abla f_k$$

The notation is usually simplified by

$$s_k = x_{k+1} - x_k$$
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Does it have a symmetric positive definite solution?

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The condition s[⊤]_k y_k > 0 is satisfied if the line search satisfies the strong Wolfe conditions.

As a corollary, we obtain the following:

Theorem 12

Assume that we use line search satisfying strong Wolfe conditions. Then in every step, the secant condition

$$\tilde{H}_{k+1}s_k = y_k$$

has a symmetric positive definite solution \tilde{H}_{k+1} .

Note that even if we demand symmetric positive definite solutions to the secant condition, there are infinitely many.

Indeed, there are n(n+1)/2 degrees of freedom in a symmetric matrix, and the secant conditions represent only *n* conditions.

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Moreover, we want to obtain \tilde{H}_{k+1} from \tilde{H}_k by

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To have a nice iterative algorithm.

We also want \tilde{H}_{k+1} to be symmetric positive definite.

We strive to choose \tilde{H}_{k+1} "close" to \tilde{H}_k .

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Consider $u = \left(y_k - \tilde{H}_k s_k\right)$

$$\tilde{H}_{k+1} = \tilde{H}_k + \frac{uu^\top}{u^\top s_k}$$

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Now, the secant condition is satisfied:

$$\tilde{H}_{k+1}s_k = \tilde{H}_k s_k + \frac{uu^{\top}s_k}{u^{\top}s_k} = \tilde{H}_k s_k + u = \tilde{H}_k s_k + \left(y_k - \tilde{H}_k s_k\right) = y_k$$

By the way, the matrix $\frac{uu^{\top}}{u^{\top}s_k}$ is of rank one and is a unique symmetric rank one matrix which makes \tilde{H}_{k+1} satisfy the secant condition.

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By the way, the matrix $\frac{uu^{\top}}{u^{\top}s_k}$ is of rank one and is a unique symmetric rank one matrix which makes \tilde{H}_{k+1} satisfy the secant condition.

To obtain a quasi-Newton method, it suffices to initialize \tilde{H}_0 , typically to the identity *I*, and use \tilde{H}_k instead of the Hessian $H_k = \nabla^2 f_k$ in Newton's method.

Symmetric Rank One Update

Algorithm 10 SR1

 $k \leftarrow 0$ $\alpha_{\text{init}} \leftarrow 1$ $H_0 \leftarrow I$ while $\|\nabla f_k\|_{\infty} > \tau$ do Solve for p_k in $\tilde{H}_k p_k = -\nabla f_k$ $\alpha \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ $x_{k+1} \leftarrow x_k + \alpha p_k$ $s \leftarrow x_{k+1} - x_k$ $y \leftarrow \nabla f_{k+1} - \nabla f_k$ $u \leftarrow v - H_k s$ $\tilde{H}_{k+1} \leftarrow \tilde{H}_k + \frac{uu^{\top}}{..^{\top}}$ $k \leftarrow k + 1$ end while

Note that the denominator $u^{\top}s_k$ can be 0, in which case the update is impossible. The usual strategy is to skip the update and set $\tilde{H}_{k+1} = \tilde{H}_k$.

We will look at a three-dimensional quadratic problem $f(x) = \frac{1}{2}x^{\top}Qx - c^{\top}x$ with

$$Q = egin{pmatrix} 2 & 0 & 0 \ 0 & 3 & 0 \ 0 & 0 & 4 \end{pmatrix}$$
 and $c = egin{pmatrix} -8 \ -9 \ -8 \end{pmatrix}$,

whose solution is $x_* = (-4, -3, -2)^{\top}$. Use the exact line search.

The initial guesses are $\tilde{H}_0 = I$ and $x_0 = (0, 0, 0)^{\top}$.

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At the initial point, $\|\nabla f(x_0)\|_{\infty} = \|-c\|_{\infty} = 9$, so this point is not optimal.

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The initial guesses are $\tilde{H}_0 = I$ and $x_0 = (0, 0, 0)^{\top}$.

At the initial point, $\|\nabla f(x_0)\|_{\infty} = \|-c\|_{\infty} = 9$, so this point is not optimal. The first search direction is

$$p_0 = \begin{pmatrix} -8 \\ -9 \\ -8 \end{pmatrix}.$$

The exact line search gives $\alpha_0 = 0.3333$.

The new estimate of the solution, the update vectors, and the new Hessian approximation are:

$$x_1 = \begin{pmatrix} -2.66 \\ -3.00 \\ -2.66 \end{pmatrix}, \nabla f_1 = \begin{pmatrix} 2.66 \\ 0 \\ -2.66 \end{pmatrix}, s_0 = \begin{pmatrix} -2.66 \\ -3.00 \\ -2.66 \end{pmatrix}, y_0 = \begin{pmatrix} -5.33 \\ -9.00 \\ -10.66 \end{pmatrix}$$

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and

$$\tilde{H}_1 = I + \frac{(y_0 - Is_0)(y_0 - Is_0)^\top}{(y_0 - Is_0)^\top s_0} = \begin{pmatrix} 1.1531 & 0.3445 & 0.4593 \\ 0.3445 & 1.7751 & 1.0335 \\ 0.4593 & 1.0335 & 2.3780 \end{pmatrix}$$

The new estimate of the solution, the update vectors, and the new Hessian approximation are:

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At this new point $\|\nabla f(x_1)\|_{\infty} = 2.66$ so we keep going, obtaining the search direction

$$p_1 = \begin{pmatrix} -2.9137 \\ -0.5557 \\ 1.9257 \end{pmatrix},$$

and the step length $\alpha_1=$ 0.3942.

This gives the new estimates:

$$x_2 = \begin{pmatrix} -3.81 \\ -3.21 \\ -1.90 \end{pmatrix}, \quad \nabla f_2 = \begin{pmatrix} 0.36 \\ -0.65 \\ 0.36 \end{pmatrix}, \quad s_1 = \begin{pmatrix} -1.14 \\ -0.21 \\ 0.75 \end{pmatrix}, \quad y_1 = \begin{pmatrix} -2.29 \\ -0.65 \\ 3.03 \end{pmatrix}$$

and

	/ 1.6568	0.6102	-0.3432	
$\tilde{H}_2 =$	0.6102	1.9153	0.6102	
	_0.3432	0.6102	3.6568 /	

At the point x_2, $\|
abla f(x_2)\|_\infty = 0.65$ so we keep going, with

$$p_2 = \begin{pmatrix} -0.4851\\ 0.5749\\ -0.2426 \end{pmatrix},$$

and $\alpha = 0.3810$.

This gives

$$x_3 = \begin{pmatrix} -4 \\ -3 \\ -2 \end{pmatrix}, \quad \nabla f_3 = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \quad s_2 = \begin{pmatrix} -0.18 \\ 0.21 \\ -0.09 \end{pmatrix}, \quad y_2 = \begin{pmatrix} -0.36 \\ 0.65 \\ -0.36 \end{pmatrix},$$

and $\tilde{H}_3 = Q$. Now $\|\nabla f(x_3)\| = 0$, so we stop.

Does symmetric rank one update satisfy our demands? We want every \tilde{H}_k to be a symmetric positive definite solution to the secant condition.

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However, for line search, let us try a bit "richer" solution to the secant condition.

Symmetric Rank Two Update

Consider

$$\tilde{H}_{k+1} = \tilde{H}_k - \frac{\left(\tilde{H}_k s_k\right) \left(\tilde{H}_k s_k\right)^\top}{s_k^\top \tilde{H}_k s_k} + \frac{y_k y_k^\top}{y_k^\top s_k}$$

Once again, verifying $\tilde{H}_{k+1}s_k = y_k$ is not difficult.

Lemma 1

Assume that \tilde{H}_k is symmetric positive definite. Then \tilde{H}_{k+1} is symmetric positive definite iff $y_k^{\top} s_k > 0$.

We know that line search satisfying the strong Wolfe conditions preserves $y_k^\top s_k > 0$.

Thus, starting with a symmetric positive definite \tilde{H}_0 (e.g., a scalar multiple of I), every \tilde{H}_k is symmetric positive definite and satisfies the secant condition.

BFGS

Algorithm 11 BFGS v1

 $k \leftarrow 0$ $\alpha_{\text{init}} \leftarrow 1$ $H_0 \leftarrow I$ while $\|\nabla f_k\|_{\infty} > \tau$ do Solve for p_k in $\tilde{H}_k p_k = -\nabla f_k$ $\alpha \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ $x_{k+1} \leftarrow x_k + \alpha p_k$ $s \leftarrow x_{k+1} - x_k$ $y \leftarrow \nabla f_{k+1} - \nabla f_k$ $\tilde{H}_{k+1} \leftarrow \tilde{H}_k - \frac{(\tilde{H}_k s)(\tilde{H}_k s)^{\top}}{s^{\top}\tilde{H}_k s} + \frac{yy^{\top}}{y^{\top}s}$ $k \leftarrow k + 1$ end while

Note that we still have to solve a linear system for p_k .

Consider the quadratic problem $f(x) = \frac{1}{2}x^{\top}Qx - c^{\top}x$ with

$$Q = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 4 \end{pmatrix} \quad \text{and} \quad c = \begin{bmatrix} -8 \\ -9 \\ -8 \end{bmatrix},$$

whose solution is $x_* = (-4, -3, -2)^{\top}$. Use the exact line search.

Choose
$$\tilde{H}_0 = I$$
 and $x_0 = (0, 0, 0)^T$.

At iteration $0, \|\nabla f(x_0)\|_{\infty} = 9$, so this point is not optimal.

The search direction is

$$p_0 = \left(\begin{array}{c} -8\\ -9\\ -8 \end{array}\right)$$

and $\alpha_0 = 0.3333$.

The new estimate of the solution and the new Hessian approximation are

$$x_1 = \begin{pmatrix} -2.6667 \\ -3.0000 \\ -2.6667 \end{pmatrix} \text{ and } \tilde{H}_1 = \begin{pmatrix} 1.1021 & 0.3445 & 0.5104 \\ 0.3445 & 1.7751 & 1.0335 \\ 0.5104 & 1.0335 & 2.3270 \end{pmatrix}$$

At iteration 1, $\|\nabla f(x_1)\|_{\infty} = 2.6667$, so we continue. The next search direction is

$$p_1 = \left(\begin{array}{c} -3.2111\\ -0.6124\\ 2.1223 \end{array}\right)$$

and $\alpha_1 = 0.3577$.

This gives the estimates.

$$x_2 = \begin{pmatrix} -3.8152 \\ -3.2191 \\ -1.9076 \end{pmatrix} \quad \text{and} \quad \tilde{H}_2 = \begin{pmatrix} 1.6393 & 0.6412 & -0.3607 \\ 0.6412 & 1.8600 & 0.6412 \\ -0.3607 & 0.6412 & 3.6393 \end{pmatrix}$$

At iteration 2, $\left\|
abla f(x_2) \right\|_\infty = 0.6572$, so we continue, computing

$$p_2 = \left(\begin{array}{c} -0.5289\\ 0.6268\\ -0.2644\end{array}\right)$$

and $\alpha_2 = 0.3495$. This gives

$$x_3 = \begin{pmatrix} -4 \\ -3 \\ -2 \end{pmatrix}$$
 and $\tilde{H}_3 = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 4 \end{pmatrix}$.

Now $\left\|\nabla f(x_3)\right\|_{\infty} = 0$, so we stop.

Notice that we got the same x_1, x_2, x_3 as for SR1. This follows from using the exact line search and the quadratic problem. It does not hold in general.


Here $\ell_1 = 12, \ell_2 = 8, k_1 = 1, k_2 = 10, mg = 7$

Two Spring Problem - BFGS



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$. Compare this with 32 iterations of gradient descent and 12 iterations of Newton's method.

Rosenbrock Function - BFGS *Rosenbrock:* $f(x_1, x_2) = (1 - x_1)^2 + 100(x_2 - x_1^2)^2$



Gradient descent, line search, stop. cond. $||\nabla f||_{\infty} \leq 10^{-6}$. Compare with 10,662 iterations of gradient descent and 24 iterations of Newton's method.

Problem: SR1 and BFGS solve $\tilde{H}_k p = -\nabla f_k$ repeatedly. What if we could iteratively update H_k^{-1} ?

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To get such a "something" we use the following Sherman–Morrison–Woodbury (SMW) formula:

$$(A + UV^{T})^{-1} = A^{-1} - A^{-1}U(I + V^{T}A^{-1}U)^{-1}V^{T}A^{-1}$$

Here A is a $(n \times n)$ -matrix, U, V are $(n \times m)$ -matrices with $m \le n$.

Rank 1 – Iterative Inverse Hessian Approximation

Applying SMW to the rank one update

$$\tilde{H}_{k+1} = \tilde{H}_k + \frac{\left(y_k - \tilde{H}_k s_k\right) \left(y_k - \tilde{H}_k s_k\right)^\top}{\left(y_k - \tilde{H}_k s_k\right)^\top s_k}$$

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yields

$$\tilde{H}_{k+1}^{-1} = \tilde{H}_{k}^{-1} + \frac{\left(s_{k} - \tilde{H}_{k}^{-1}y_{k}\right)\left(s_{k} - \tilde{H}_{k}^{-1}y_{k}\right)^{\top}}{\left(s_{k} - \tilde{H}_{k}^{-1}y_{k}\right)^{\top}y_{k}}$$

Yes, only y and s swapped places.

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Yes, only y and s swapped places.

This allows us to avoid solving $\tilde{H}_k p_k = -\nabla f_k$ for p_k in every iteration.

Rank One Update V2

Algorithm 12 Rank 1 update v1 1: $k \leftarrow 0$ 2: $\alpha_{\text{init}} \leftarrow 1$ 3: $H_0 \leftarrow I$ 4: while $\|\nabla f_k\|_{\infty} > \tau$ do $p_k \leftarrow -\tilde{H}_{l_k}^{-1} \nabla f_k$ 5: 6: $\alpha \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ 7: $x_{k+1} \leftarrow x_k + \alpha p_k$ 8: $s \leftarrow x_k - x_{k-1}$ 9: $\mathbf{v} \leftarrow \nabla f_k - \nabla f_{k-1}$ $\tilde{H}_{k+1}^{-1} \leftarrow \tilde{H}_{k}^{-1} + \frac{(s - \tilde{H}_{k}^{-1} y) (s - \tilde{H}_{k}^{-1} y)^{\top}}{(s - \tilde{H}_{k}^{-1} y)^{\top} y}$ 10: $k \leftarrow k + 1$ 11: 12: end while

BFGS

Applying SMW to the BFGS Hessian update

$$\tilde{H}_{k+1} = \tilde{H}_k - \frac{\left(\tilde{H}_k s_k\right) \left(\tilde{H}_k s_k\right)^\top}{s_k^\top \tilde{H}_k s_k} + \frac{y_k y_k^\top}{y_k^\top s_k}$$

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Applying SMW to the BFGS Hessian update

$$\tilde{H}_{k+1} = \tilde{H}_k - \frac{\left(\tilde{H}_k s_k\right) \left(\tilde{H}_k s_k\right)^\top}{s_k^\top \tilde{H}_k s_k} + \frac{y_k y_k^\top}{y_k^\top s_k}$$

yields

$$\tilde{H}_{k+1}^{-1} = \left(I - \frac{s_k y_k^\top}{s_k^\top y_k}\right) \tilde{H}_k^{-1} \left(I - \frac{y_k s_k^\top}{s_k^\top y_k}\right) + \frac{s_k s_k^\top}{s_k^\top y_k}$$

We avoid solving the linear system for p_k .

BFGS V2

Algorithm 13 BFGS v2 1: $k \leftarrow 0$ 2: $\alpha_{\text{init}} \leftarrow 1$ 3: $H_0 \leftarrow I$ 4: while $\|\nabla f_k\|_{\infty} > \tau$ do $p_k \leftarrow -\tilde{H}_k^{-1} \nabla f_k$ 5: 6: $\alpha \leftarrow \text{linesearch}(p_k, \alpha_{\text{init}})$ 7: $x_{k+1} \leftarrow x_k + \alpha p_k$ 8: $k \leftarrow k+1$ 9: $s \leftarrow x_k - x_{k-1}$ 10: $y \leftarrow \nabla f_k - \nabla f_{k-1}$ $\tilde{H}_{k+1}^{-1} \leftarrow \left(I - \frac{sy^{\top}}{s^{\top}y}\right) \tilde{H}_{k}^{-1} \left(I - \frac{ys^{\top}}{s^{\top}y}\right) + \frac{ss^{\top}}{s^{\top}y}$ 11: 12: end while

Let us denote by s_0, \ldots, s_k and y_0, \ldots, y_k the values of the variables s and y, resp., during the iterations $1, \ldots, k$ of BFGS.

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So, the matrix \tilde{H}_k does not have to be stored if the algorithm remembers the values s_0, \ldots, s_k and y_0, \ldots, y_k .

Note that this would be more space efficient for k < n.

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This is the basic idea behind limited-memory BFGS which stores only the running window s_{k-m}, \ldots, s_k and y_{k-m}, \ldots, y_k and computes \tilde{H}_k using these values as if initialized by $\tilde{H}_{k-m-1} = I$.

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The space complexity becomes nm, which is beneficial when n is large.

Another View on BFGS (Optional)

We search for \tilde{H}_{k+1}^{-1} where \tilde{H}_{k+1} satisfies $\tilde{H}_{k+1}s_k = y_k$. Search for a solution \tilde{V} for $\tilde{V}y_k = s_k$.

The idea is to use \tilde{V} close to \tilde{H}_k^{-1} (in some sense):

$$\min_{ ilde{H}} \left\| ilde{V} - ilde{H}_k^{-1}
ight\|$$

subject to $ilde{V} = ilde{V}^ op, \quad ilde{V} y_k = s_k$

Here the norm is weighted Frobenius norm:

$$\|A\| \equiv \left\| W^{1/2} A W^{1/2} \right\|_F,$$

where $\|\cdot\|_F$ is defined by $\|C\|_F^2 = \sum_{i=1}^n \sum_{j=1}^n c_{ij}^2$. The weight W can be chosen as any matrix satisfying the relation $Wy_k = s_k$.

BFGS is obtained with $W = \overline{G}_k^{-1}$ where \overline{G}_k is the average Hessian defined by $\overline{G}_k = \left[\int_0^1 \nabla^2 f(x_k + \tau \alpha_k p_k) d\tau\right]$

Solving this gives precisely the BFGS formula for \tilde{H}_{k+1}^{-1} .

Global Convergence of Line Search

Denote by θ_k the angle between p_k and $-\nabla f_k$, i.e., satisfying

$$\cos \theta_k = \frac{-\nabla f_k^T p_k}{\|\nabla f_k\| \, \|p_k\|}$$

Recall that f is L-smooth for some L > 0 if

$$\|
abla f(x) -
abla f(ilde{x})\| \le L \|x - ilde{x}\|, \quad \text{ for all } x, ilde{x} \in \mathbb{R}^n$$

Theorem 13 (Zoutendijk)

Consider $x_{k+1} = x_k + \alpha_k p_k$, where p_k is a descent direction and α_k satisfies the strong Wolfe conditions. Suppose that f is bounded below, continuously differentiable, and L-smooth. Then

$$\sum_{k\geq 0}\cos^2\theta_k\,\|\nabla f_k\|^2<\infty.$$

Global Convergence of Quasi-Newton's Method

Assume that all α_k satisfy strong Wolfe conditions.

Assume that the approximations to the Hessians \tilde{H}_k are positive definite with a uniformly bounded condition number:

$$\left|\left| ilde{H}_k
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Then θ_k between $p_k = -\tilde{H}_k^{-1} \nabla f_k$ and $-\nabla f_k$ and satisfies

 $\cos \theta_k \ge 1/M$

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 for all k

Then θ_k between $p_k = -\tilde{H}_k^{-1} \nabla f_k$ and $-\nabla f_k$ and satisfies

 $\cos \theta_k \ge 1/M$

Thus, under the assumptions of Zoutendijk's theorem, we obtain

$$\frac{1}{M^{2}} \sum_{k \ge 0} \|\nabla f_{k}\|^{2} \le \sum_{k \ge 0} \cos^{2} \theta_{k} \|\nabla f_{k}\|^{2} < \infty$$

which implies that $\lim_{k\to\infty} ||\nabla f_k|| = 0$.

Behavior of BFGS

▶ It may happen that \tilde{H}_k becomes a poor approximation of the Hessian H_k . If, e.g., y_k^{\top} is tiny, then \tilde{H}_{k+1} will be huge.

However, it has been proven experimentally that if \tilde{H}_k wrongly estimates the curvature of f and this estimate slows down the iteration, then the approximation will tend to correct the bad Hessian approximations.

The above self-correction works only if an appropriate line search is performed (strong Wolfe conditions).

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The above self-correction works only if an appropriate line search is performed (strong Wolfe conditions).

There are more sophisticated ways of setting the initial Hessian approximation H₀.

See Numerical Optimization, Nocedal & Wright, page 201.

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- Each iteration is performed for O(n²) operations as opposed to O(n³) for methods involving solutions of linear systems.
- There is even a memory-limited variant (L-BFGS) that uses only information from past *m* steps, and its single iteration complexity is O(*mn*).

- Each iteration is performed for O(n²) operations as opposed to O(n³) for methods involving solutions of linear systems.
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- There is even a memory-limited variant (L-BFGS) that uses only information from past *m* steps, and its single iteration complexity is O(*mn*).
- Compared with Newton's method, no second derivatives are computed.
- Local superlinear convergence can be proved under specific conditions.

Compare with local quadratic convergence of Newton's method and linear convergence of gradient descent.

Limited-Memory BFGS

When the number of design variables is extensive, working with the whole Hessian inverse approximation matrix might not be practical.

This motivates limited-memory quasi-Newton methods,

In addition, these methods also improve the computational efficiency of medium-sized problems (hundreds or thousands of design variables) with minimal sacrifice in accuracy.

L-BFGS

Recall that we compute iteratively the approximation to the inverse Hessian by

$$H_{k+1}^{-1} = \left(I - \frac{s_k y_k^\top}{s_k^\top y_k}\right) H_k^{-1} \left(I - \frac{y_k s_k^\top}{s_k^\top y_k}\right) + \frac{s_k s_k^\top}{s_k^\top y_k}$$

However, eventually, we are interested in

$$p_k = H_k^{-1} \nabla f$$

Note that given the sequences s_1, \ldots, s_k and y_1, \ldots, y_k and H_0^{-1} we can recursively compute H_{k+1}^{-1} for every k.

What if we limit the sequences in memory to just m last elements:

$$s_{k-m+1}, s_{k-m+2}, \dots, s_k$$
 $y_{k-m+1}, y_{k-m+2}, \dots, y_k$

In practice, m between 5 and 20 is usually sufficient. We also initialize the recurrence with the last iterate:

L-BFGS

Let us rewrite the BFGS update formula as follows:

$$\tilde{H}_{k+1}^{-1} = V_k^T \tilde{H}_k^{-1} V_k + \rho_k s_k s_k^\top$$

where

$$\begin{aligned} \rho_k &= s_k^\top y_k \quad \text{and} \quad V_k &= I - \rho_k s_k y_k^\top \\ s_k &= x_{k+1} - x_k \quad \text{and} \quad y_k &= \nabla f_{k+1} - \nabla f_k \end{aligned}$$

By substitution, we obtain

$$\begin{split} \tilde{H}_{k}^{-1} &= \left(V_{k-1}^{T} \cdots V_{k-m}^{T} \right) \tilde{H}_{k}^{0} \left(V_{k-m} \cdots V_{k-1} \right) \\ &+ \rho_{k-m} \left(V_{k-1}^{T} \cdots V_{k-m+1}^{T} \right) s_{k-m} s_{k-m}^{T} \left(V_{k-m+1} \cdots V_{k-1} \right) \\ &+ \rho_{k-m+1} \left(V_{k-1}^{T} \cdots V_{k-m+2}^{T} \right) s_{k-m+1} s_{k-m+1}^{T} \left(V_{k-m+2} \cdots V_{k} \right) \\ &+ \cdots \\ &+ \cdots \\ &+ \rho_{k-1} s_{k-1} s_{k-1}^{T} \end{split}$$

L-BFGS Algorithm

Algorithm 14 L-BFGS two-loop recursion **Input:** : $s_{k-1}, ..., s_{k-m}$ and $y_{k-1}, ..., y_{k-m}$ **Output:** : p_k the search direction $-\tilde{H}_k^{-1}\nabla f_k$ 1: $a \leftarrow \nabla f_{k}$ 2: for $i = k - 1, k - 2, \dots, k - m$ do 3: $\alpha_i \leftarrow \rho_i s_i^T q$ 4: $q \leftarrow q - \alpha_i y_i$ 5: end for 6: $r \leftarrow H^0_{\mu} q$ 7: for $i = k - m, k - m + 1, \dots, k - 1$ do 8: $\beta \leftarrow \rho_i v_i^T r$ 9: $r \leftarrow r + s_i(\alpha_i - \beta)$ 10: end for 11: stop with result $\tilde{H}_{\iota}^{-1} \nabla f_k = r$
L-BFGS Algorithm

Algorithm 15 L-BFGS

- 1: Choose starting point x_0 , integer m > 0
- 2: $k \leftarrow 0$
- 3: repeat

4: Choose
$$H_k^0$$
 e.g. $\frac{s_{k-1}^{-1}y_{k-1}}{y_{k-1}^{-1}y_{k-1}}$

- 5: Compute $p_k \leftarrow -H_k \nabla f_k$ using the previous algorithm
- 6: Compute $x_{k+1} \leftarrow x_k + \alpha_k p_k$, where α_k is chosen to satisfy the strong Wolfe conditions
- 7: **if** k > m **then**
- 8: Discard the vector pair $\{s_{k-m}, y_{k-m}\}$ from storage
- 9: end if
- 10: Compute and save $s_k \leftarrow x_{k+1} x_k$, $y_k \leftarrow \nabla f_{k+1} \nabla f_k$
- 11: $k \leftarrow k+1$

12: **until** convergence

$$f(x_1, x_2) = (1 - x_1)^2 + (1 - x_2)^2 + \frac{1}{2} (2x_2 - x_1^2)^2$$

Stopping: $||\nabla f||_{\infty} \leq 10^{-6}$.



In L-BFGS, the memory length m was 5. The results are similar.



Here $\ell_1 = 12, \ell_2 = 8, k_1 = 1, k_2 = 10, mg = 7$



Rosenbrock: $f(x_1, x_2) = (1 - x_1)^2 + 100(x_2 - x_1^2)^2$



Rosenbrock:

$$f(x_1, x_2) = (1 - x_1)^2 + 100(x_2 - x_1^2)^2$$



Computational Complexity

Algorithm	Computational Complexity
Steepest Descent	O(n) per iteration
Newton's Method	$O(n^3)$ to compute Hessian and solve system
BFGS	$O(n^2)$ to update Hessian approximation

Table: Summary of the computational complexity for each optimization algorithm.

- Steepest Descent: Simple but often slow, requiring many iterations.
- Newton's Method: Fast convergence but expensive per iteration.
- BFGS: Quasi-Newton, no Hessian needed, good speed and iteration count balance.