

IV054 Coding, Cryptography and Cryptographic Protocols  
2014 - Exercises VI.

1. Using Fermat's little theorem, determine an integer  $x$  such that  $8^x \equiv 2 \pmod{23}$ .
2. How many cars do you have to observe in order for the probability to be greater than 50% of observing at least two cars with the same first three symbols on their license plate? A car license plate consists of 7 symbols: there is a capital letter (A, ..., Z) in the second position, other positions are occupied with digits. Each symbol is equally probable.
3. Let  $p$  be an odd prime number and  $g$  be a primitive root modulo  $p$ . This means that the powers  $1, g, g^2, \dots, g^{p-2}$  are all distinct modulo  $p$ . Suppose  $m$  is an odd number. Prove that  $g^m$  is a quadratic nonresidue modulo  $p$ .
4. Suppose we know how to factorize large numbers and we find out that the private keys of the Rabin cryptosystem are  $p = 163$  and  $q = 307$ . Decrypt the cryptotext "15244 33337". The message is a meaningful English word.

5. Let  $f, g$  be negligible functions. Prove the following:

- (a)  $f^k$  is negligible for any  $k > 0, k \in \mathbb{R}$ .
- (b)  $f + g$  is negligible.

(You can use the alternative definition of negligible function: it is enough for the function to be smaller than  $n^{-c}$  (for every positive integer  $c$ ) rather than  $\frac{1}{p(n)}$ .)

6. Show that you can factorize efficiently if you have an oracle for finding square roots. Demonstrate by factorizing  $n = 88416763$  in case the oracle tells you that the square roots of  $51733469 \pmod{88416763}$  are 50224876, 38191887, 22222, 88394541.
7. Using the Shank's algorithm find  $x$  such that

$$88^x = 80 \pmod{107}$$

and show all steps of the algorithm.

8. Consider the Rabin cryptosystem with  $n = p^k q^s$  where  $k, s > 1$ .
  - (a) How many possible plaintexts do we obtain after decryption?
  - (b) Find a decryption of cryptotext  $c = 14590$  using private keys  $11^{27^3} = 41503$ .

Hint: (Hensel's Lemma)

If  $r_i$  is square root of  $a \pmod{p^i}$  then

$$r_{i+1} = r_i + tp^i$$

is a square root of  $a \pmod{p^{i+1}}$  where  $t$  is solution of

$$t2r_i \equiv -C \pmod{p}$$

and  $C = \frac{r_i^2 - a}{p^i}$  (here always  $p^i | r_i^2 - a$ ).