

Limits of sequences of Latin squares

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- Graphs**: Borgs-Chayes-Lovász-Sós-Szegedy-Vesztergombi (2007+)
- Permutations**: Hoppen-Kohayakawa-Moreira-Ráth-Sampaio (2013)

Limits of permutations

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$X = \{x_1 < \dots < x_m\} \subseteq [k]$ we have

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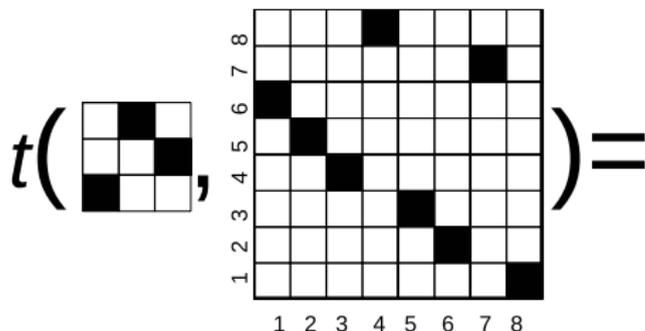
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$$\tau = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}, \sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 6 & 5 & 4 & 8 & 3 & 2 & 7 & 1 \end{pmatrix}$$



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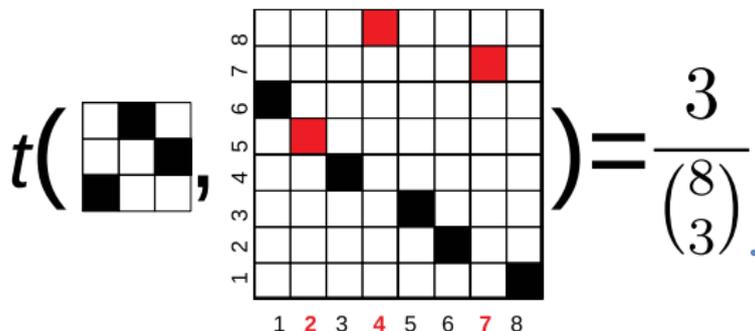
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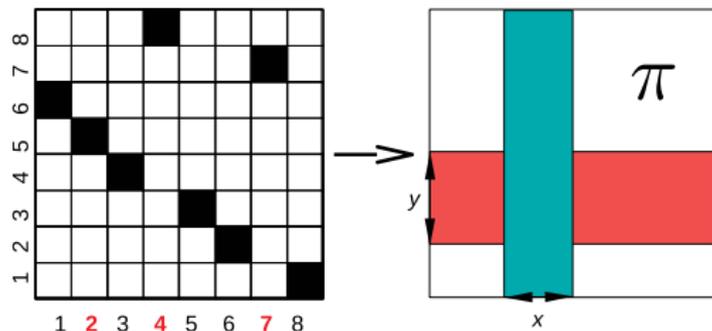

$$t\left(\begin{array}{|c|c|c|} \hline \blacksquare & \blacksquare & \square \\ \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array}, \begin{array}{|c|c|c|c|c|c|c|} \hline \square & \square & \square & \color{red}\blacksquare & \square & \square & \square & \square \\ \hline \blacksquare & \square & \square & \square & \square & \square & \color{red}\blacksquare & \square \\ \hline \square & \color{red}\blacksquare & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \blacksquare & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \blacksquare & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \blacksquare & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \blacksquare & \square & \square \\ \hline \square & \blacksquare \\ \hline \end{array}\right) = \frac{3}{\binom{8}{3}}.$$

Limits of permutations

A permuton P is a probability measure on $[0, 1]^2$ that has uniform marginals, i.e. $P(A \times [0, 1]) = P([0, 1] \times A) = \lambda(A)$, where λ is the Lebesgue measure.

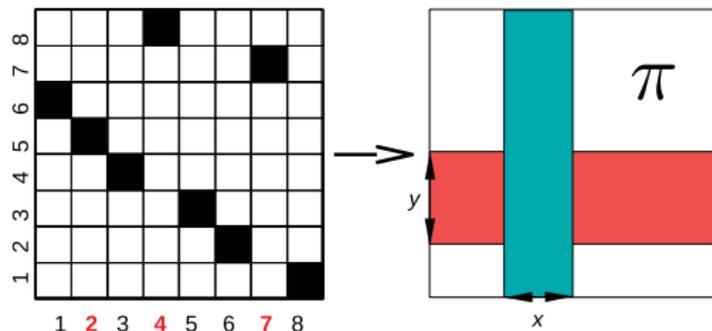
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For this theory a compactness result, a cut-norm, and counting lemmas were developed by Hoppen, Kohayakawa, Moreira, Ráth and Sampaio ('13, JCTB).

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One can think of a Latin square as a **2-dimensional permutation**.

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$$L[r_i, c_j] \leq L[r_{i'}, c_{j'}] \text{ iff } A_{i,j} \leq A_{i',j'} \quad \forall i, i' \in [k], j, j' \in [\ell].$$

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Example

$$A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix} \quad L = \begin{matrix} 1 & 5 & 3 & 6 & 4 & 2 \\ 4 & 3 & 6 & 1 & 2 & 5 \\ 3 & 4 & 5 & 2 & 1 & 6 \\ 6 & 2 & 4 & 3 & 5 & 1 \\ 2 & 6 & 1 & 5 & 3 & 4 \\ 5 & 1 & 2 & 4 & 6 & 3 \end{matrix}$$

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Limit objects - Motivational examples

Standard Cyclic Example

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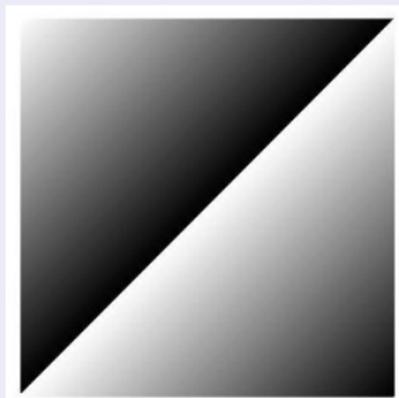
0	1	0	1	2	0	1	2	3	0	1	2	3	4
1	0	1	2	0	1	2	3	0	1	2	3	4	0
		2	0	1	2	3	0	1	2	3	4	0	1
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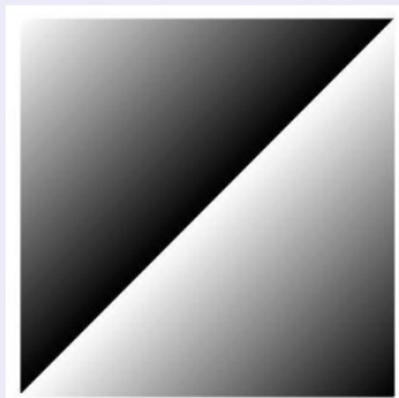


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So should the limit object
be a function

$$L : [0, 1]^2 \rightarrow [0, 1] ?$$

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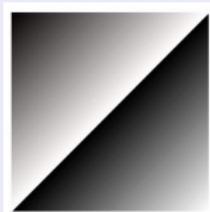
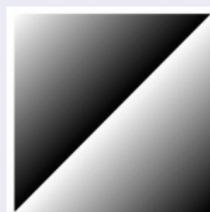
Probabilistic example

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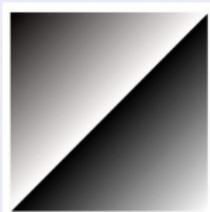
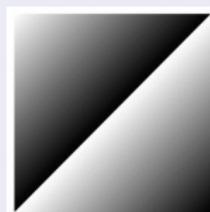
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\mathcal{P} is the space of probability distributions on $[0, 1]$ and

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Limit objects - Motivational examples

Odd-even example

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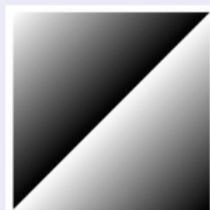
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Equivalently: (ν_W, f) , where f is as above and ν_W is a probability measure on $\Omega^2 \times [0, 1]$ with uniform marginals related to the above definition via

$$\nu_W(S \times T \times V) = \int_S \int_T W(x, y)(V) dx dy .$$

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Limit of the cyclic example

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$\Omega = [0, 1]$, f is the identity and $W : [0, 1]^2 \rightarrow \mathcal{P}$ is defined by

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Limit of the odd-even example

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$$\Omega = [0, 1] \times \{\text{odd}, \text{even}\},$$

$$f : [0, 1] \times \{\text{odd}, \text{even}\} \rightarrow [0, 1], (x, a) \mapsto x$$

$W : ([0, 1] \times \{\text{odd}, \text{even}\})^2 \rightarrow \mathcal{P}$ is defined by

$$W((x, a), (y, b)) := \begin{cases} \text{Dirac}(x + y \pmod 1) & \text{if } a = b, \\ \text{Dirac}(-x - y \pmod 1) & \text{if } a \neq b. \end{cases}$$

Densities in Latinons

For a Latinon $L = (W, f)$, define

$t(A, L) := \mathbb{P}(A^* \equiv A \mid \text{when a } k \times \ell \text{ matrix } A^* \text{ is 'sampled' from } (W, f)) .$

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$t(A, L) := \mathbb{P}(A^* \equiv A \mid \text{when a } k \times \ell \text{ matrix } A^* \text{ is 'sampled' from } (W, f))$.

Sampling: Select $(x_1, \dots, x_k) \in \Omega^k$ and $(y_1, \dots, y_\ell) \in \Omega^\ell$ with

$f(x_1) < f(x_2) < \dots < f(x_k)$ and $f(y_1) < f(y_2) < \dots < f(y_\ell)$ u.a.r.

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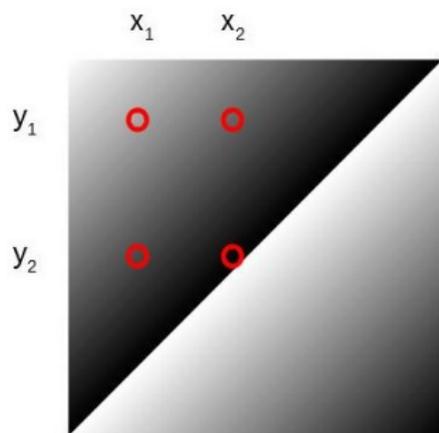
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Example

For $A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}$,

$t(A, (W, f)) = \frac{1}{6}$.

Densities in Latinons

For a matrix $A \in \mathbb{R}^{k \times \ell}$ we set

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Densities in Latinons

Let (W, f) be a Latinon and $A \in \mathbb{R}^{k \times \ell}$. We denote by $t(A, (W, f))$ the *density* of the *pattern* A in (W, f) and define it to be

$$t(A, (W, f)) := k! \ell! \int_{\mathbf{x} \in [0, 1]_{<f}^k} \int_{\mathbf{y} \in [0, 1]_{<f}^\ell} \left(\bigotimes_{(i,j) \in [k] \times [\ell]} W(x_i, y_j) \right) (\mathcal{R}^A([0, 1])) dy dx .$$

Limit theories of discrete structures

- (F1) **Finite discrete structures** and substructures together with a notion of density $t(\cdot, \cdot)$. (E.g. homomorphism density in graphs.)
- (F2) **Left-convergence**: A sequence of structures (S_n) is left-convergent if $(t(H, S_n))$ converges for every finite substructure H .
- (F3) **Limit objects**: We define a space of analytic limits objects and the notion of density is extended to those limit objects. (Graphons.)
- (F4) **Compactness**: Every sequence of structures contains a subsequence converging to a limit object. (Lovász-Szegedy '06.)
- (F5) **Denseness**: For every limit object there exists a converging sequence of discrete structures. (W -random graphs.)
- (F6) **Equivalence of local and global**: There is another 'global' metric generating the same topology as left-convergence. (Cut-distance.)

Compactness theorem

Let $(L_n)_{n \in \mathbb{N}}$ be a sequence of Latin squares or Latinons, and let (W, f) be a Latinon. We say $L_n \rightarrow (W, f)$ if $\lim_{n \rightarrow \infty} t(A, L_n) = t(A, (W, f))$ for every $k, \ell \in \mathbb{N}$ and $k \times \ell$ pattern A .

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Compactness for Latinons (G., Hancock, Hladký, Sharifzadeh, 20⁺)

Let $(L_n)_{n \in \mathbb{N}}$ be a sequence of Latinons. There exists a subsequence $(L_{n_i})_{i \in \mathbb{N}}$ and a Latinon (W, f) such that

$$L_{n_i} \rightarrow (W, f).$$

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Continuous image of compact space is compact, hence $\iota^{-1}(\iota(\mathcal{L})) = \mathcal{L}$ is compact.

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- Define compression $\iota((W, f)) := (O^f, W_{1,1}, W_{1,2}, \dots)$.

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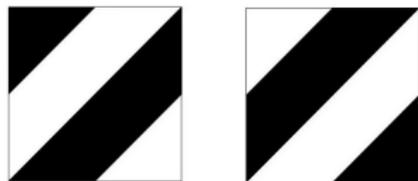
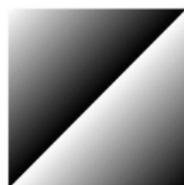


Figure: $W_{1,1}$ and $W_{1,2}$

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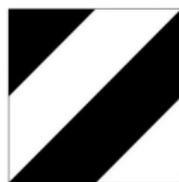


Figure: $W_{2,1}, W_{2,2}, W_{2,3}, W_{2,4}$

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Cut-distance for Latinons

Cut-distance for graphons W and U

$\delta_{\square}(W, U) := \inf_{\varphi \in \mathcal{S}_{[0,1]}} \|W - U^{\varphi}\|_{\square}$ where

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Cut-distance for Latinons $L_1 = (W, f)$ and $L_2 = (U, g)$

$$\delta_L(L_1, L_2) := \inf_{\varphi, \psi \in \mathcal{S}_{\Omega}} (\|W - U^{\varphi, \psi}\|_L + \|O^f - O^{g \circ \varphi}\|_{\square} + \|O^f - O^{g \circ \psi}\|_{\square})$$

where $O : \Omega^2 \rightarrow [0, 1]$ is a graphon s.t. $O(x, y) := \begin{cases} 1, & x < y, \\ 0, & \text{otherwise;} \end{cases}$

$$\|W - U^{\varphi, \psi}\|_L := \sup_{\substack{R, C \subseteq \Omega, \\ V \subseteq [0,1] \text{ interval}}} \left| \int_{x \in R} \int_{y \in C} W(x, y)(V) - U(\varphi(x), \psi(y))(V) dy dx \right|.$$

Motivation for the cut-distance

$$L_n(i, j) := \begin{cases} i + j \pmod n & \text{if } i + j \equiv 0 \pmod 2, \\ -i - j \pmod n & \text{if } i + j \equiv 1 \pmod 2. \end{cases}$$

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0	5	2	3	4	1	0	1	4	3	2	5
5	2	3	4	1	0	1	4	3	2	5	0
2	3	4	1	0	5	4	3	2	5	0	1
3	4	1	0	5	2	3	2	5	0	1	4
4	1	0	5	2	3	2	5	0	1	4	3
1	0	5	2	3	4	5	0	1	4	3	2

Equivalence of local and global

Counting Lemma (G., Hancock, Hladký, Sharifzadeh, 20⁺)

Let $k, \ell \in \mathbb{N}$. Then there exists a constant $c_{k,\ell}$ such that for every $d \in \mathbb{N}$, Latinons L_1, L_2 and $k \times \ell$ pattern A we have

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Equivalence

Convergence w.r.t. densities $t(\cdot, \cdot) \iff$ convergence w.r.t. cut-distance δ_L .

Minimality

Approximation (G., Hancock, Hladký, Sharifzadeh, 20⁺)

For each Latinon (W, f) there exists a sequence $(L_n)_{n \in \mathbb{N}}$ of finite Latin squares of growing orders such that

$$L_n \rightarrow (W, f) .$$

Proof idea - Rödl nibble + Keevash

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- 4 Use tools from Keevash's theory about designs to extend the approximate triangle decomposition (partial Latin square) to a triangle decomposition (complete Latin square).

Further Questions

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