

Structure and Generation of Crossing-critical Graphs

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joint work with

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Definition. Graph *H* is *c*-crossing-critical if $cr(H) \ge c$ and cr(H - e) < c for all edges $e \in E(H)$.

We study crossing-critical graphs to understand what structural properties force the crossing number of a graph to be large.

Some starting examples

• Kuratowski (30): The only 1-crossing-critical graphs K_5 and $K_{3,3}$.





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 Širáň (84), Kochol (87): Infinitely many *c*-crossing-critical graphs for every *c* ≥ 2, even simple 3-connected.



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• Hliněný (02):



... and a bit of surprise

 Dvořák, Mohar (10): A *c*-crossing-crit. graph with unbounded degree, *c* ≥ 171.



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- Bokal, Oporowski, Richter, Salazar (16): Fully described 2-crossing-critical graphs up to fin. small exceptions.
- Dvořák, Hliněný, Mohar, Postle (11, not published):
 A *c*-crossing-critical graph cannot contain a deep nest, and so it has bounded dual diameter.

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+ combinations of these together

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- II. There are well-defined local operations (replacements) that can reduce any large *c*-crossing-critical graph to a smaller one.

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* Our result *

- I. "Nothing else than the previous" can constitute crossing-criticality.
- II. There are well-defined local operations (replacements) that can reduce any large *c*-crossing-critical graph to a smaller one.
- III. There are finitely many well-defined building bricks that can produce all *c*-crossing-critical graphs from a finite set of base graphs.

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- **1.** General understanding of the struct. of a plane band and tiles:
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 - or, get a topological long-band structure composed of boundedsize tiles separated (between consecutive ones) by paths.



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Prove that such shortening preserves crossing-criticality.



Starting from a path-decomposition of bounded width, the main trouble is that its bags do not correspond to our topological graph (our picture).



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 Apply an algebraic tool Simon's factorization forest, to a semigroup formed by concatenation of these topological types.
- **c)** The previous gives a subband with a "homogeneous topol. structure"; either the desired band with properly separated and connected tiles, or one of special substructures forbidden in crossing-critical graphs:







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c) Use further repetitions of this local picture around to argue that *c*-crossing-criticality is preserved:

- G_1 drawn with < c crossings \rightarrow can expand with no new crossing,
- (more difficult) G e drawn with < c crossings \rightarrow can modify and shrink to $G_1 e$ with no new crossing.

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Thank you for your attention.