IV054 Coding, Cryptography and Cryptographic Protocols 2019 - Exercises VI.

- 1. (4 points)
 - (a) Encrypt your UČO using the Rabin cryptosystem with n = 698069. Then calculate all four possible decryptions of the ciphertext you calculated, with the knowledge that $n = 887 \times 787$.
 - (b) Encrypt your UČO with the ElGamal cryptosystem with p = 567899, q = 2, x = 12345 and random choice r = 938.
- 2. (2 points) Using baby-step giant-step algorithm find x such that $7^x \equiv 505 \pmod{541}$.
- 3. (4 points)
 - (a) What is the probability that at least two students attending IV054 this year have the same birthday?
 - (b) What is the probability that another student of IV054 is sharing birthday with you?

(135 students attend IV054 in 2019.)

- 4. (3 points) Consider a cryptographic hash function with 64 bits long output. Provide an approximate number of random guesses you have to perform to find a collision with probability at least 3/4.
- 5. (2 points) Consider the Blum-Goldwasser cryptosystem with parameters p = 11 and q = 43. Encode the message x = 1111 with $s_0 = 195$.
- 6. (4 points) Determine whether the following functions are negligible. Prove your statement.
 - (a) $\ln\left(1+\frac{1}{n}\right)$
 - (b) $e^{1/n}e^{-n}$
- 7. (5 points)
 - (a) Suppose you know a valid plaintext-ciphertext pair $w_1 = 457$, $(a_1, b_1) = (663, 2138)$, constructed using the ElGamal cryptosystem with public key p = 6661, q = 6, y = 6015. Also you know that instead of using a new random r to encrypt each new message, the sender just increments the previous one, i.e. $r_2 = r_1 + 1$. With this knowledge decrypt the following ciphertext $(a_2, b_2) =$ (3978, 1466) without calculating discrete logarithms.
 - (b) Show that the same attack is possible for any linear update function of the random seed, i.e. whenever $r_2 = kr_1 + \ell \mod p 1$.
- 8. (6 points) Consider the following cryptosystem. Let n = pq where p and q are primes. The value n is made public, $(n, \phi(n))$ forms private key.

To encrypt a message $m \in \mathbb{Z}_n$, choose a random $r \in \mathbb{Z}_n^*$ and compute

$$c = (1+n)^m r^n \mod n^2.$$

Decryption: To decrypt a ciphertext c, compute

$$m = \frac{(c^{\phi(n)} \mod n^2) - 1}{n} \cdot \phi(n)^{-1} \mod n$$

where integer division is used.

- (a) Let n = 3953. Use r = 1111 to encrypt m = 2019. Decrypt the obtained ciphertext using the fact $n = 59 \times 67$.
- (b) Decrypt c = 4354044 without using the private key, only with the knowledge of the plaintextciphertext pair from (a).
- (c) Prove the following fact that is exploited in this cryptosystem: For integers n and $a, 1 \le a < n$: $(1+n)^a = 1 + an \pmod{n^2}$.