## IV054 Coding, Cryptography and Cryptographic Protocols 2017 - Exercises X.

- 1. Consider the coin-flipping by telephone protocol (*Protocol 2* from the lecture). Let p = 11, q = 23 and x = 52. Show computation steps in detail.
- 2. Alice and Bob have a boolean function  $f : \{0,1\} \times \{0,1\} \rightarrow \{0,1\}$  known to both of them. Show that if f is not a constant function there is no protocol that lets them perfectly evaluate f(a,b) on all respective secret inputs a and b without neither party obtaining any information about the other party's secret.
- 3. Consider 30 combination padlocks where Victor knows the combinations for all of them. Give a protocol that allows Peggy to prove to Victor that she knows the combination for at least one of the padlocks without revealing which one.
- 4. Consider the following commitment scheme, with the following public information: p a large prime, g a generator of  $\mathbb{Z}_p^*$  and  $h = g^k \mod p$  with 0 < k < p-1 a random integer not known to any party. The commitment function is

 $\operatorname{commit}(r, x) = (g^r \mod p, h^{r+x} \mod p),$ 

where x is the committed bit and 0 < r < p a random integer.

- (a) Find a reveal phase for this protocol.
- (b) Discuss the binding and hiding properties of this protocol. Are they computationally/information-theoretically secure?
- (c) What happens if Bob (the receiver) knows  $\log_q(h)$ ?
- 5. Alice and Bob are using the *Bit commitment scheme II* from the lecture slides, but Bob does not trust Alice and wants her to commit her bit b twice using the same  $\alpha$  but different  $x_0$  and  $x_1$  to create commitments  $f_0 = f(b, x_0)$  and  $f_1 = f(b, x_1)$ . Bob is also lazy and instead of opening both of the commitments separately and checking if his received commitments  $f_0$  and  $f_1$  are equal to  $f(b, x_0)$  and  $f(b, x_1)$ , respectively, he just checks whether  $f(b, x_0 + x_1) = f_0 \cdot f_1$ . Will this protocol open the bit commitment correctly?
- 6. Show that the following variants of oblivious transfer are equivalent:
  - Rabin oblivious transfer: Alice transmits a bit b to Bob, who receives either b or  $\perp$  (indicating that the bit was not received), each case with probability  $\frac{1}{2}$ . Alice does not know which is the case.
  - 1-out-of-2 oblivious transfer: Alice has two bits  $b_0$  and  $b_1$ . Bob chooses  $c \in \{0, 1\}$ , learns  $b_c$  but not  $b_{1-c}$ . Alice does not learn c.
  - 1-out-of-k oblivious transfer: Alice has k bits  $b_1, \ldots, b_k$ . Bob chooses  $c \in \{1, \ldots, k\}$ , learns  $b_c$  but none of the others. Alice does not learn c.
- 7. Victor is color-blind and cannot distinguish between colors at all. Peggy who can see colors has two apples, one green and one red, but otherwise identical. Design a zero-knowledge protocol that allows Peggy to convince Victor that the apples have different colors.

More on next page >>>

## $8. \ Bonus \ exercise$

This is the title page of a book called *Steganographia*.



You can find the copy of this picture in the study materials.