## 2012 - Exercises I.

1. Let $d \geq d^{\prime}, q, n \in \mathbb{N}$. Show that $A_{q}\left(n, d^{\prime}\right) \geq A_{q}(n, d)$.
2. Compute in detail code rate of the following binary codes:
(a) $C_{1}=\{1000,1101,0010,0111\}$,
(b) $C_{2}=\{1000,0110,0111,1101,0010\}$,
(c) a code with $n=7$ and $M=16$.
3. Every bit sent through a binary erasure channel is substituted with a special symbol $e$ ( $e$ stands for erasure) with probability $p$.
(a) Suppose $n$-bit codewords are used. How many different $n$-symbol strings can appear as the channel output?
(b) For the binary erasure channel, derive an upper bound analogous to the sphere packing bound.
4. How many valid ISBN numbers start with $08493852 x_{9} x_{10}$ ? Which of these books do you consider useful for study of cryptography?
5. (a) Construct a Huffman code for the letters A, B, C, D, E and F with the frequency of use given below.

| Letter | Frequency |
| :---: | :--- |
| A | $30 \%$ |
| B | $22 \%$ |
| C | $15 \%$ |
| D | $13 \%$ |
| E | $10 \%$ |
| F | $10 \%$ |

(b) Find the average word length of the proposed Huffman code.
6. Suppose $C$ is a binary optimal prefix code for a set of messages $S=\left\{x_{1}, \ldots, x_{n}\right\}$, where message $x_{i}$ occurs with probability $p_{i}$. Prove the following statements.
(a) If $p_{j}>p_{k}$, then $l_{j} \leq l_{k}$, where $l_{i}$ is the length of codeword for the message $x_{i}$.
(b) The two longest codewords have the same length.
(c) The two longest codewords differ only in the last bit and correspond to the two least likely messages

A prefix code is a code such that no codeword is a prefix of any other codeword.
An optimal prefix code $C$ is a prefix code with minimal average length, ie. if $C^{\prime}$ is another prefix code for $S$ then

$$
\sum_{i=1}^{n} l_{i} p_{i} \leq \sum_{i=1}^{n} l_{i}^{\prime} p_{i}
$$

where $l_{i}^{\prime}$ is the length of codeword for the message $x_{i}$ in $C^{\prime}$.

