IV054 Coding, Cryptography and Cryptographic Protocols 2012 - Exercises I.

- 1. Let $d \ge d'$, $q, n \in \mathbb{N}$. Show that $A_q(n, d') \ge A_q(n, d)$.
- 2. Compute in detail code rate of the following binary codes:
 - (a) $C_1 = \{1000, 1101, 0010, 0111\},\$
 - (b) $C_2 = \{1000, 0110, 0111, 1101, 0010\},\$
 - (c) a code with n = 7 and M = 16.
- 3. Every bit sent through a binary *erasure* channel is substituted with a special symbol e (e stands for erasure) with probability p.
 - (a) Suppose *n*-bit codewords are used. How many different *n*-symbol strings can appear as the channel output?
 - (b) For the binary erasure channel, derive an upper bound analogous to the sphere packing bound.
- 4. How many valid ISBN numbers start with $08493852x_9x_{10}$? Which of these books do you consider useful for study of cryptography?
- 5. (a) Construct a Huffman code for the letters A, B, C, D, E and F with the frequency of use given below.

Letter	Frequency
А	30%
В	22%
\mathbf{C}	15%
D	13%
Ε	10%
\mathbf{F}	10%

- (b) Find the average word length of the proposed Huffman code.
- 6. Suppose C is a binary optimal prefix code for a set of messages $S = \{x_1, \ldots, x_n\}$, where message x_i occurs with probability p_i . Prove the following statements.
 - (a) If $p_j > p_k$, then $l_j \leq l_k$, where l_i is the length of codeword for the message x_i .
 - (b) The two longest codewords have the same length.
 - (c) The two longest codewords differ only in the last bit and correspond to the two least likely messages

A prefix code is a code such that no codeword is a prefix of any other codeword.

An optimal prefix code C is a prefix code with minimal average length, *ie.* if C' is another prefix code for S then

$$\sum_{i=1}^{n} l_i p_i \le \sum_{i=1}^{n} l'_i p_i$$

where l'_i is the length of codeword for the message x_i in C'.